LABORATOIRE D'INFORMATIQUE DE NANTES ATLANTIQUE
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Learning from an informant

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PASCAL
Pattern Analysis, Statistical Modelling and Computational Learning

cnrs
dépasser les frontières

Région
PAYS DE LA LOIRE

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- List is necessarily incomplete. Excuses to those that have been forgotten.

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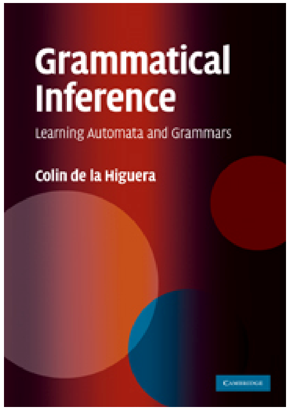
Chapter 12 (and part of 14)



Outline

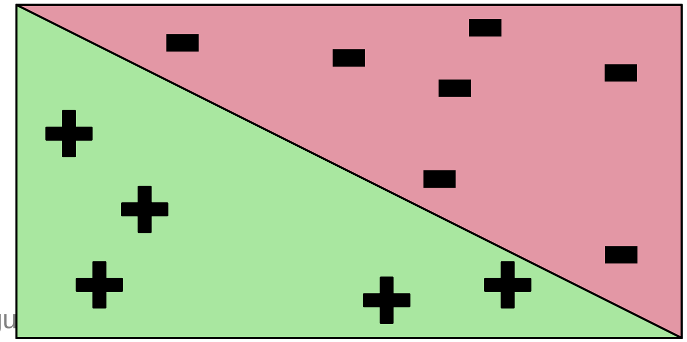
1. The rules of the game
2. Basic elements for learning DFA
3. Gold's algorithm
4. RPNI
5. Complexity discussion
6. Heuristics
7. Open questions and conclusion

1 The rules of the game

Motivation

- We are given a set of strings S_+ and a set of strings S_-
- Goal is to build a classifier
- This is a traditional (or typical) machine learning question
- How should we solve it?

Σ^*

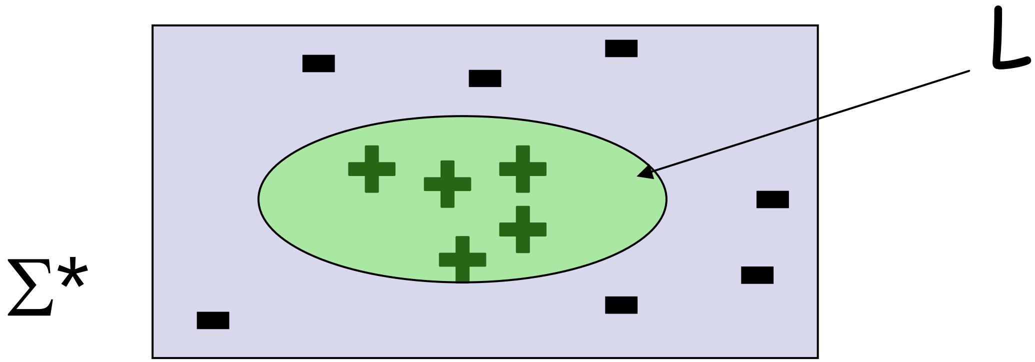


Ideas

- Use a distance between strings and try k-NN
- Embed strings into vectors and use some off-the-shelf technique (decision trees, SVMs, other kernel methods)

Alternative

- Suppose the classifier is some grammatical formalism
- Thus we have L and Σ*\L



Informed presentations

- An *informed* presentation (or an informant) of $L \subseteq \Sigma^*$ is a function $f : \mathbb{N} \rightarrow \Sigma^* \times \{-,+\}$ such that $f(\mathbb{N})=(L,+)\cup(\overline{L},-)$
- f is an infinite succession of all the elements of Σ^* labelled to indicate if they belong or not to L.

Obviously many possible candidates

- Any Grammar G such that
  - $S_+ \subseteq \mathbf{L}(G)$
  - $S_- \cap \mathbf{L}(G) = \emptyset$
- But there is an infinity of such grammars!

Structural completeness

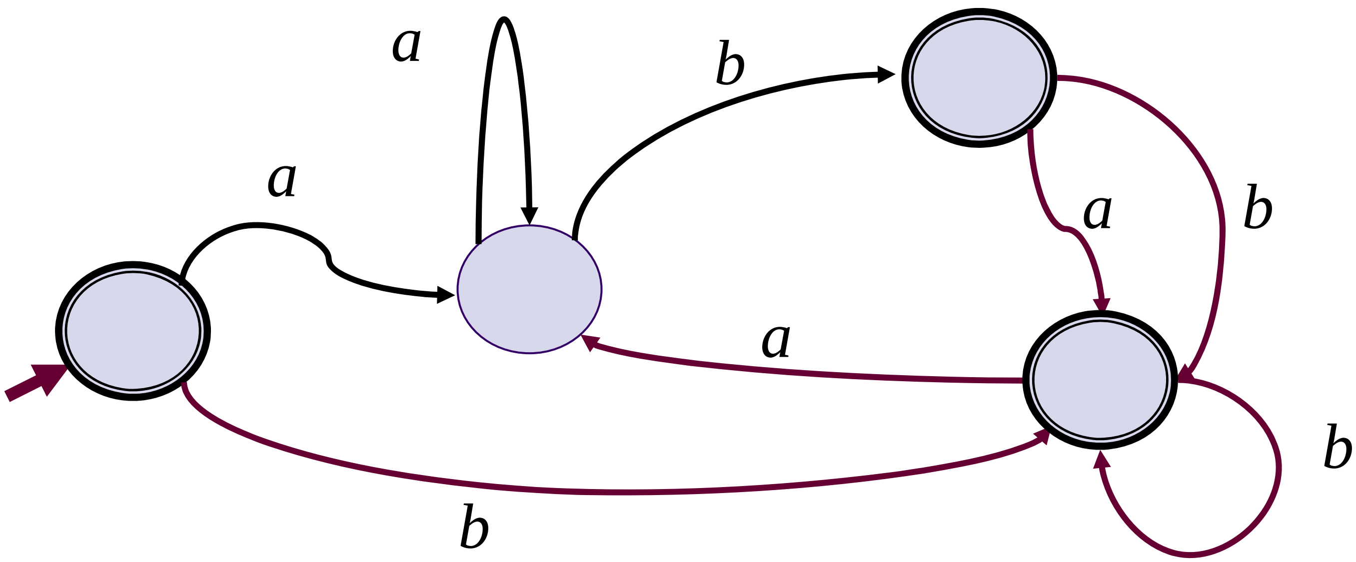
- (of S_+ *re* a *DFA*)

  each edge is used at least once
  each final state accepts at least one string

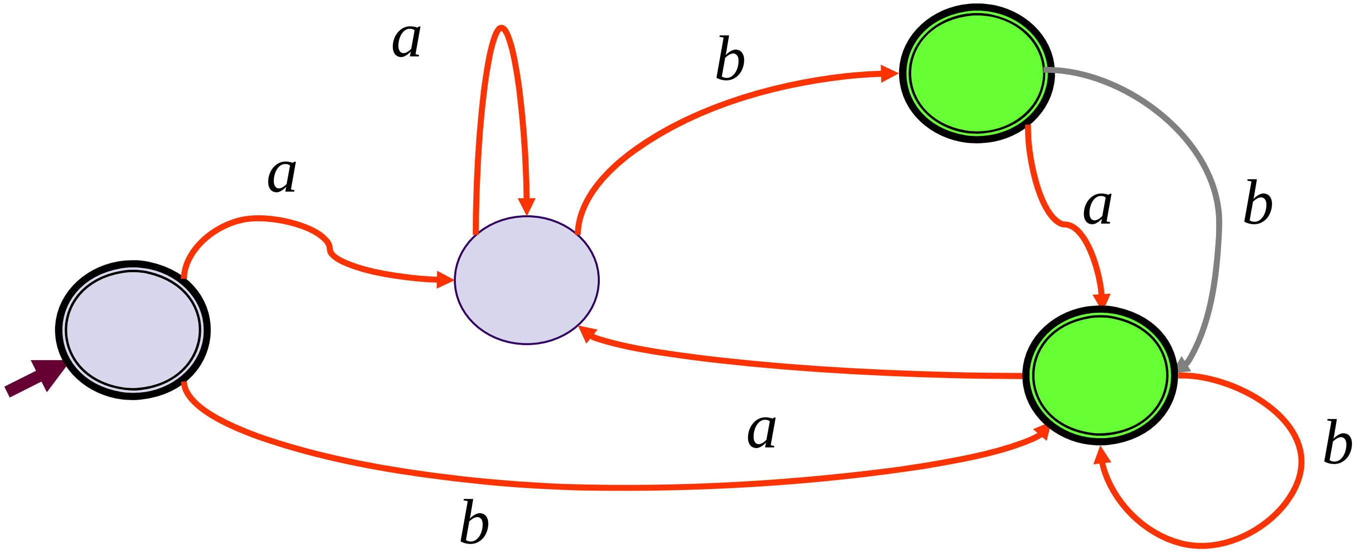
- Look only at *DFA* for which the sample is structurally complete!

Example

- S_+={*aab*, *b*, *aaaba*, *bbaba*}



- S_+={*aab*, *b*, *aaaba*, *bbaba*} …

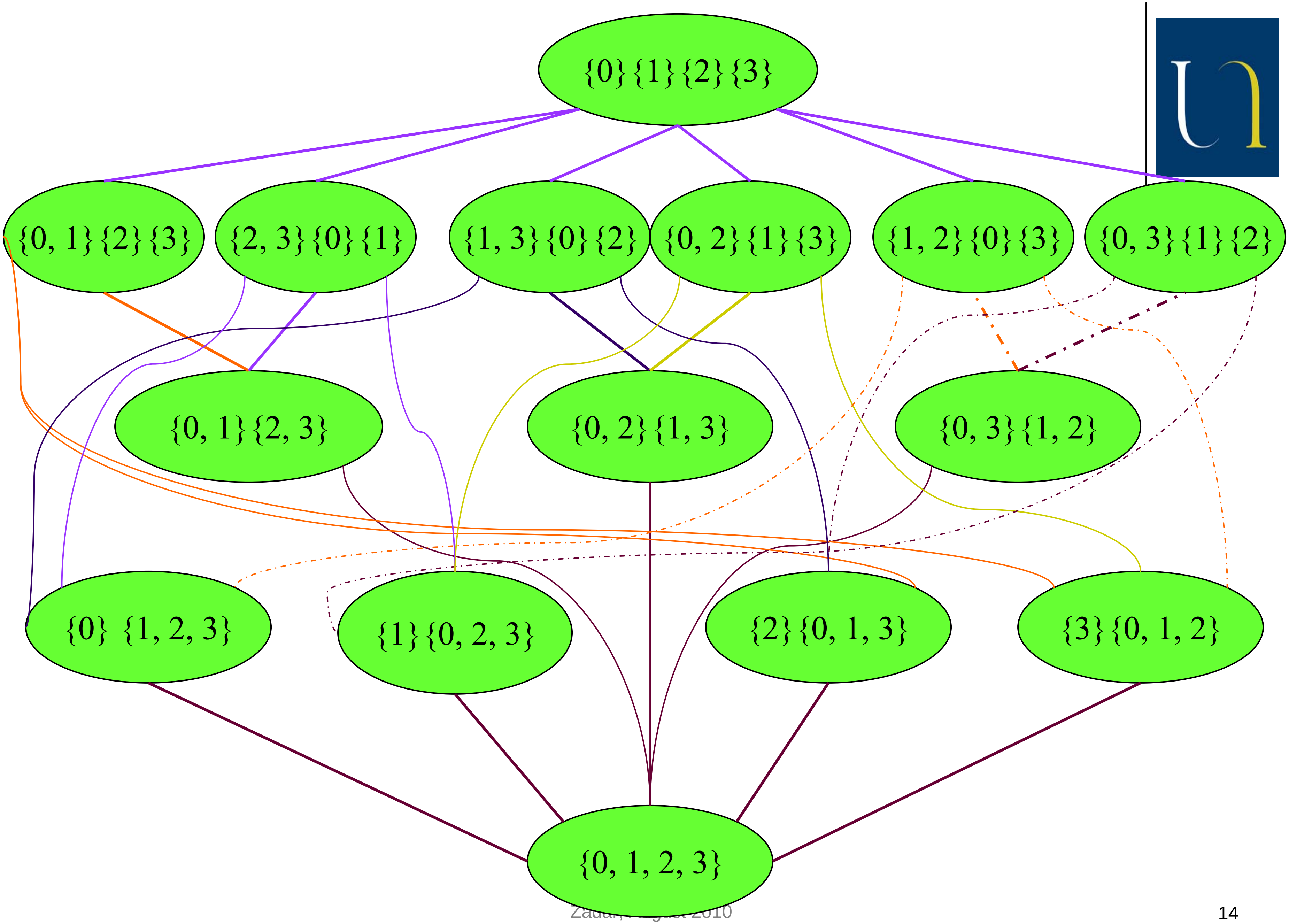


add λ and *abb*

Defining the search space by structural completeness

(Dupont, Miclet, Vidal 94)

- the basic operation: merging two states
- a bias on the concepts : structural completeness of the positive sample S_+
- a theorem: every biased solution can and can only be obtained by merging states in $CA(S_+)$
- the search space is a partition lattice.


{0} {1} {2} {3}
{0, 1} {2} {3}
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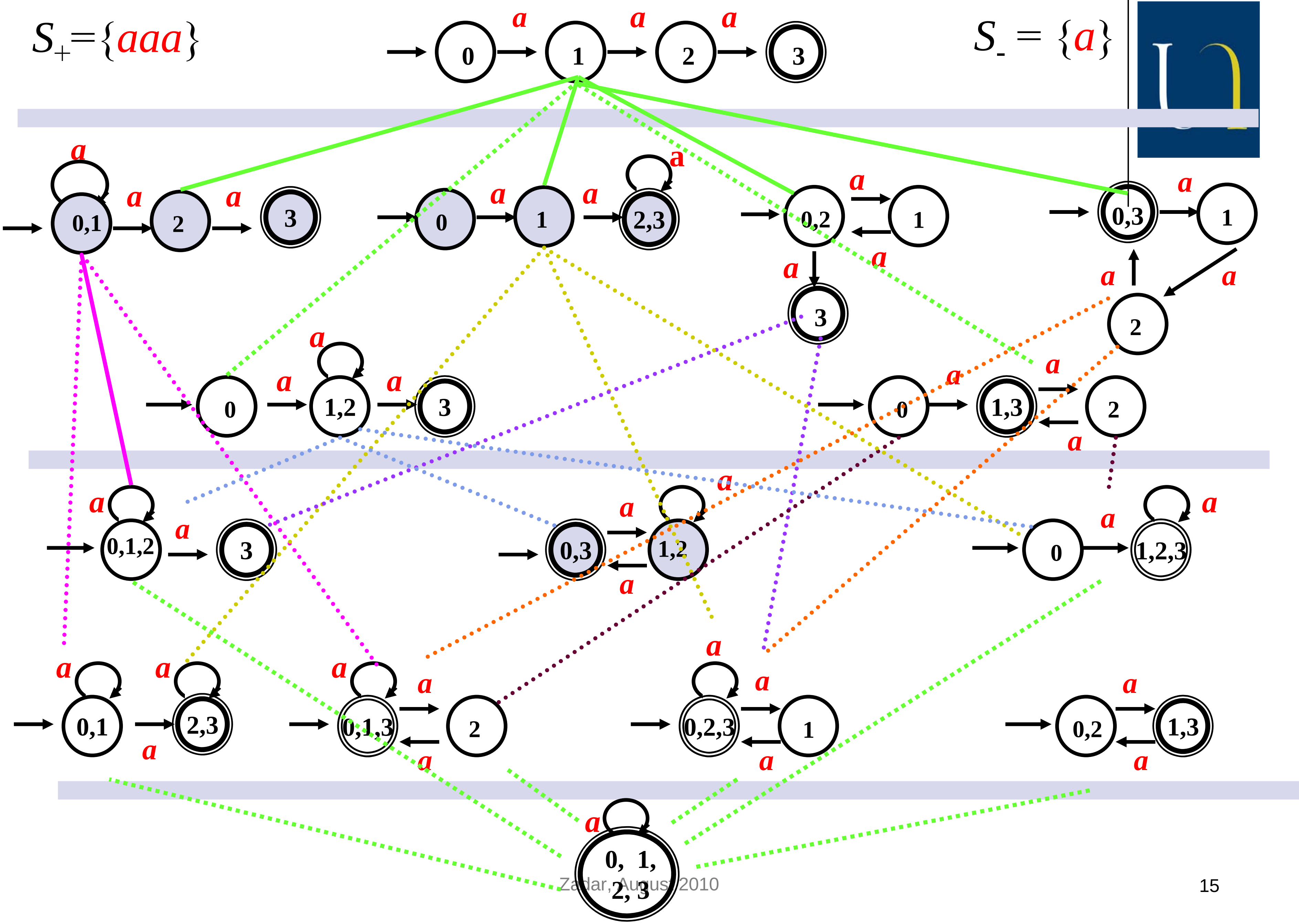
{0, 3} {1, 2}
{0} {1, 2, 3}
{1} {0, 2, 3}
{2} {0, 1, 3}
{3} {0, 1, 2}
{0, 1, 2, 3}

$S_+ = \{aaa\}$

$S_- = \{a\}$



The partition lattice

- Let E be a set with n elements
- The number of partitions of E is given by the Bell number

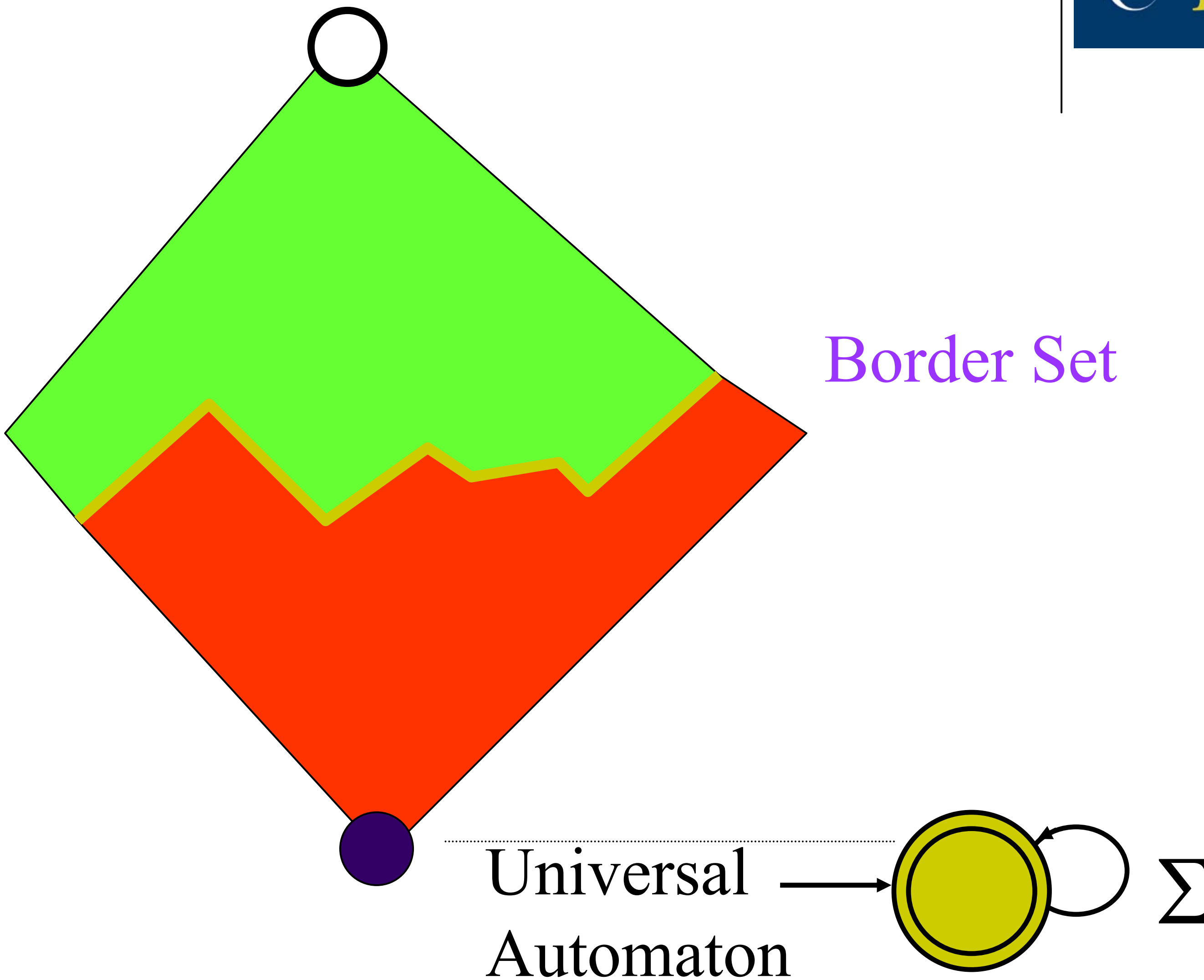
$$\begin{cases} \omega(0) = 1 \\ \omega(n+1) = \sum_{p=0}^{n} C_n^p . \omega(n) \end{cases}$$

$\omega(16)=$10 480 142 147

Regular inference as search

- another result: the smallest *DFA* fitting the examples is in the lattice constructed on $PTA(S_+)$
- generally, algorithms would start from $PTA(S_+)$ and explore the corresponding lattice of solutions using the merging operation. S_- is used to control the generalization.

$CA\ (S_+)$ or $PTA(S_+)$

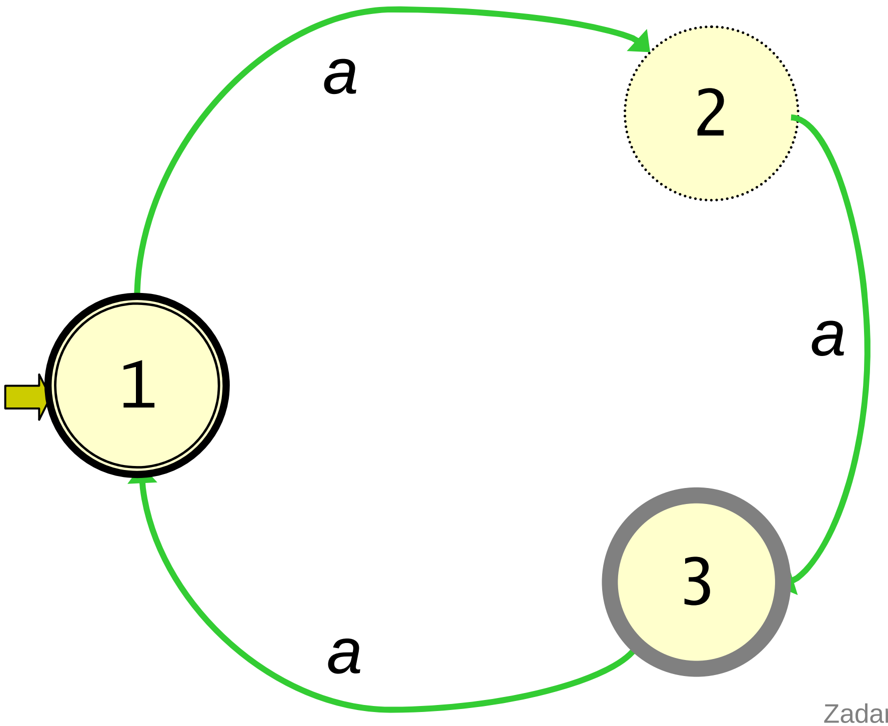


2 Basic structures

Two types of final states

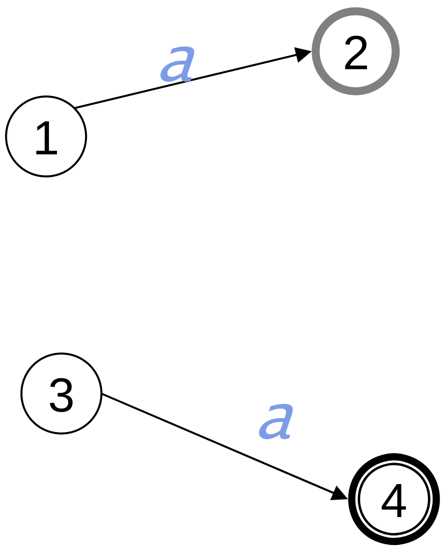
$S_+=\{\lambda,\ aaa\}$

$S_-=\{aa,\ aaaaa\}$

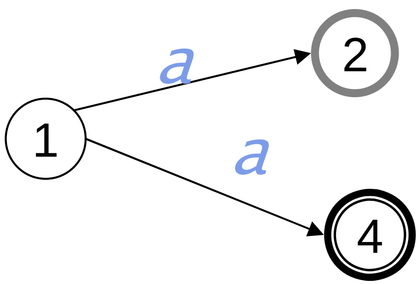


1 is accepting
3 is rejecting
What about state 2?

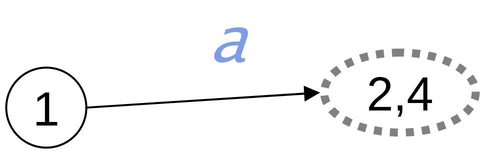
What is determinism about?



Merge 1 and 3?



But…

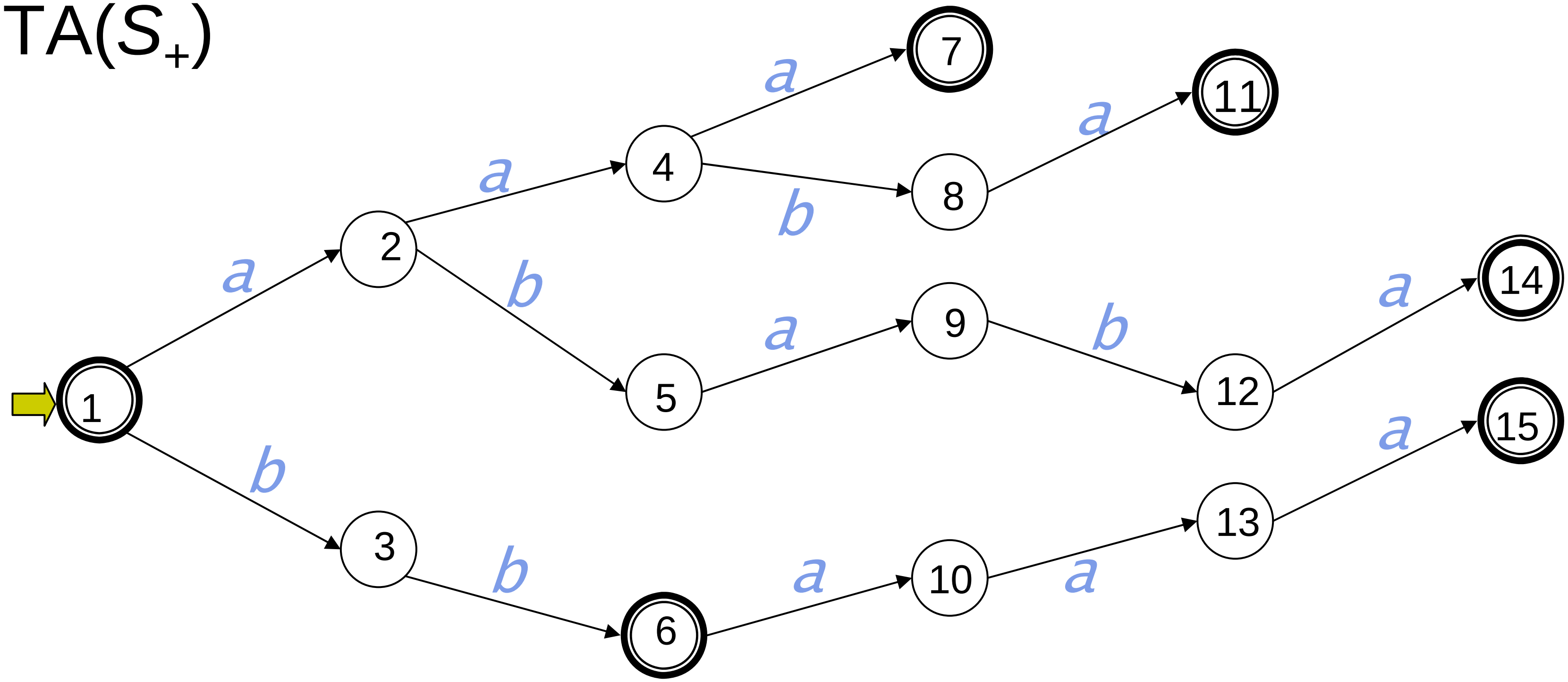


The prefix tree acceptor

- The smallest tree-like DFA consistent with the data
- Is a solution to the learning problem
- Corresponds to a rote learner

From the sample to the PTA

PTA(S_+)

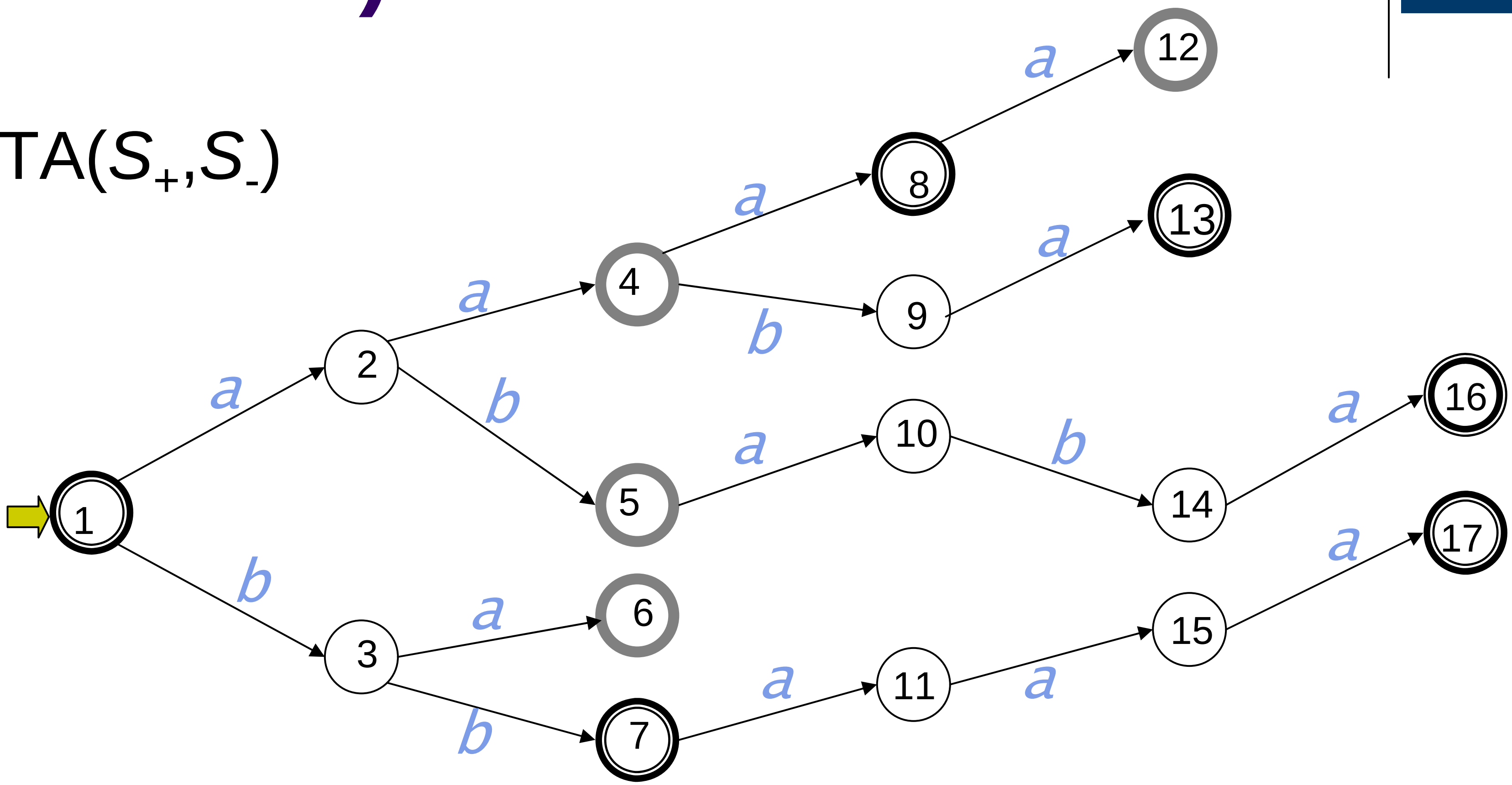


$S_+=\{\lambda,\ aaa,\ aaba,\ ababa,\ bb,\ bbaaa\}$

$S_-=\{aa,\ ab,\ aaaa,\ ba\}$

From the sample to the PTA (full PTA)

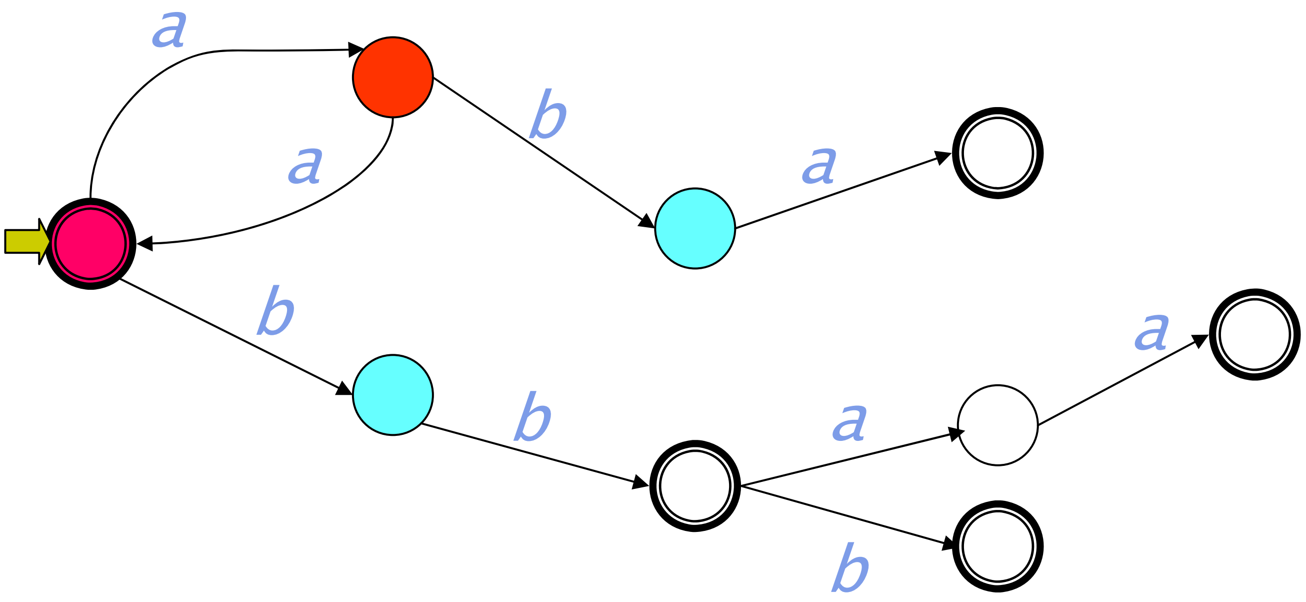
PTA(S_+,S_-)



$S_+=\{\lambda,\ aaa,\ aaba,\ ababa,\ bb,\ bbaaa\}$
$S_-=\{aa,\ ab,\ aaaa,\ ba\}$

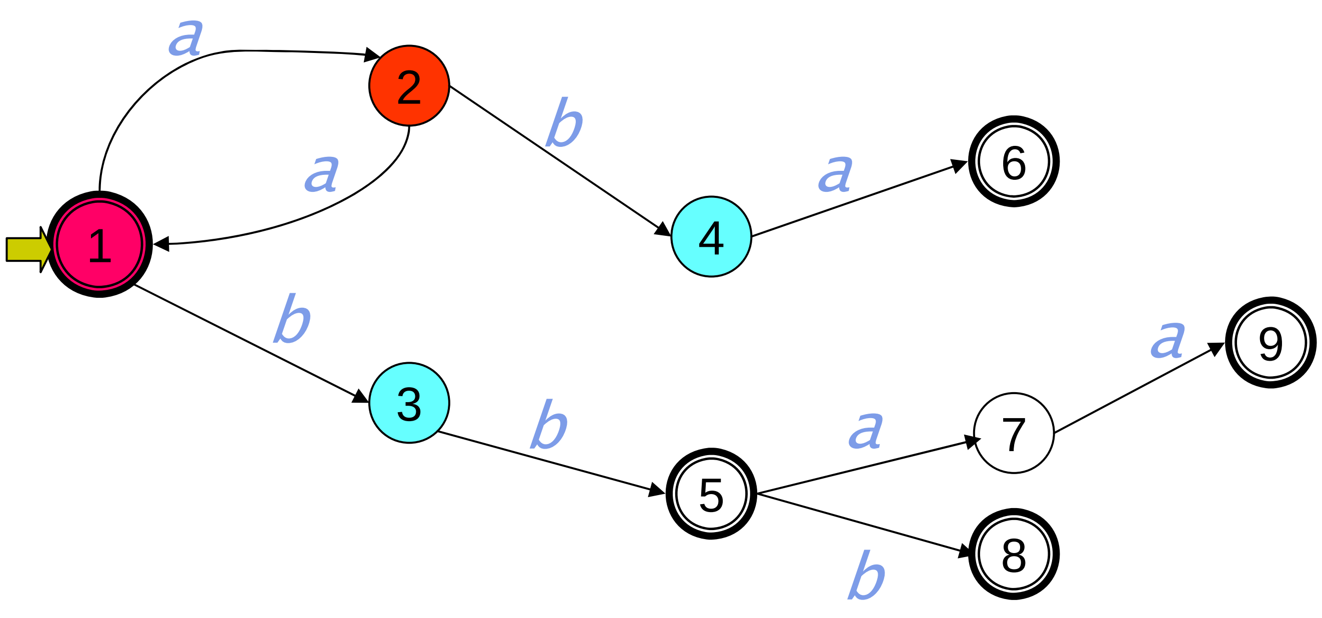
Red, Blue and White states

-Red states are confirmed states
-Blue states are the (non Red) successors of the Red states
-White states are the others



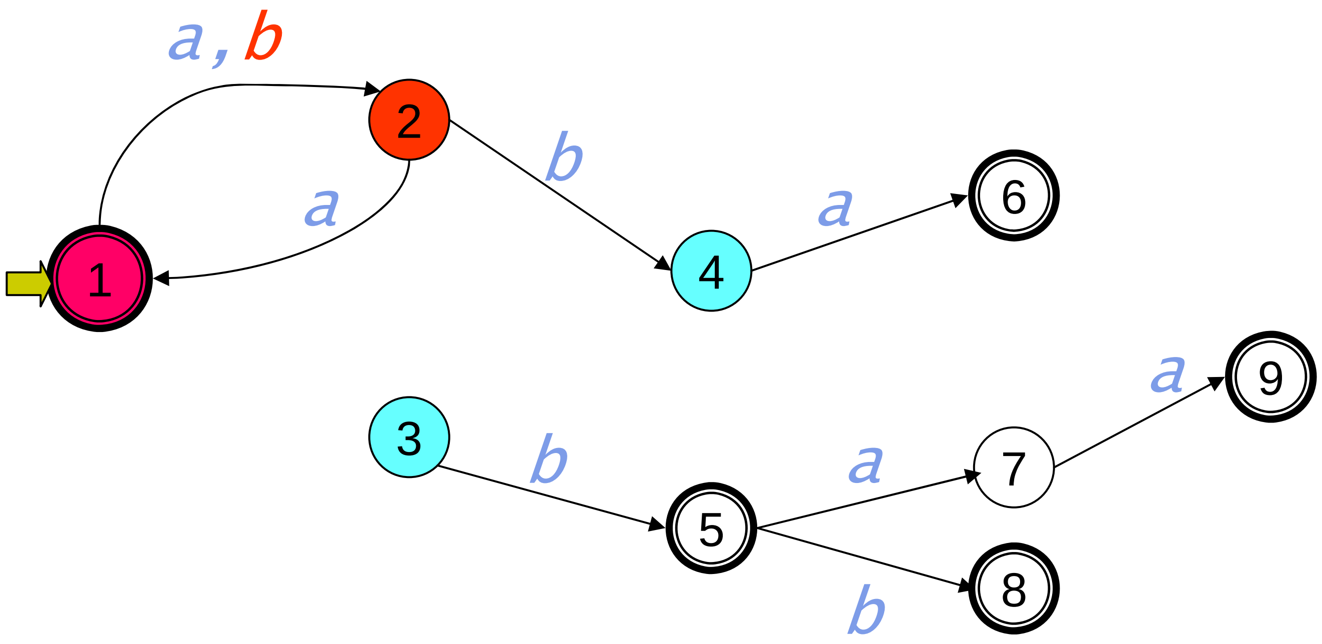
Merge and fold

Suppose we want to merge state 3 with state 2



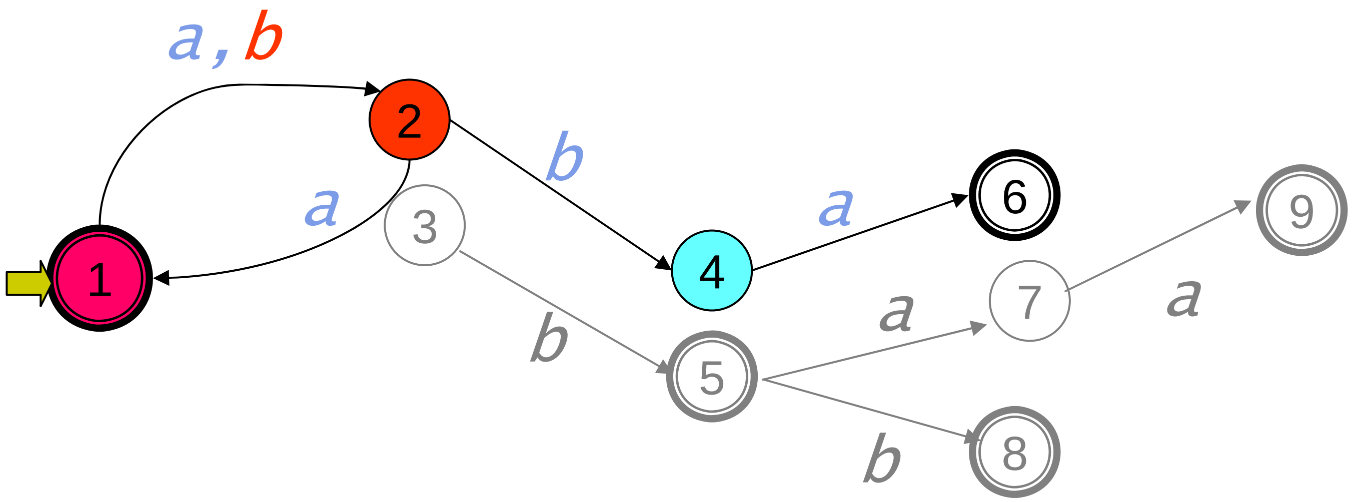
Merge and fold

First disconnect 3
and reconnect to 2



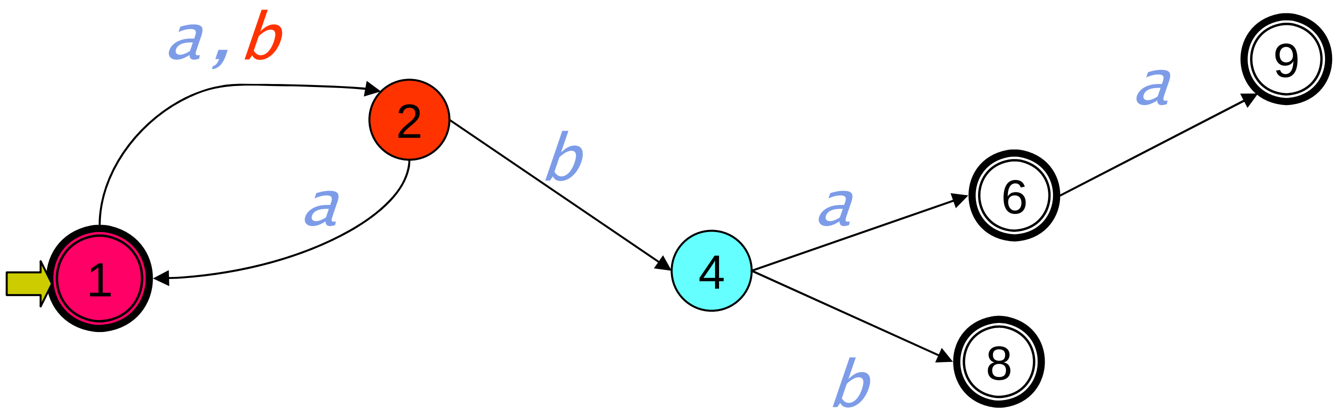
Merge and fold

Then fold subtree rooted in 3 into the DFA starting in 2



Merge and fold

Then fold subtree rooted in 3 into the DFA starting in 2



Other search spaces

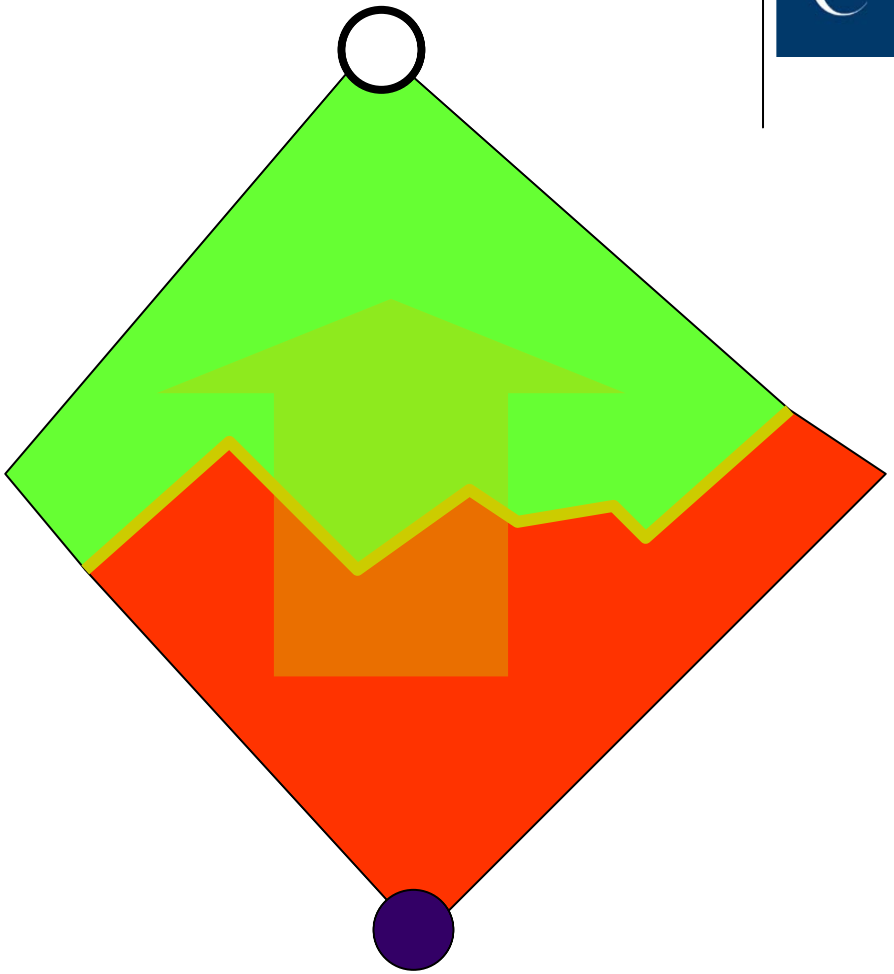
an augmented *PTA* can be constructed on both S_+ and S_- (Coste 98, Oliveira 98)

- but not every merge is possible
- the search algorithms must run under a set of dynamic constraints

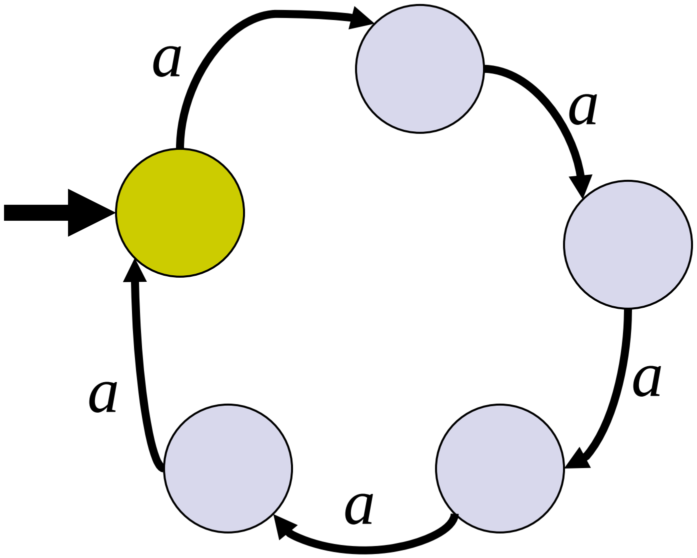
State splitting

Searching by splitting

- start from the one-state universal automaton, keep constructing *DFA* controlling the search with $\langle S_+, S_- \rangle$



- *That seems a good idea… but take a^{5*}. What 4 (or 3, 2, 1) state automaton is a decent approximation of a^{5*} ?*



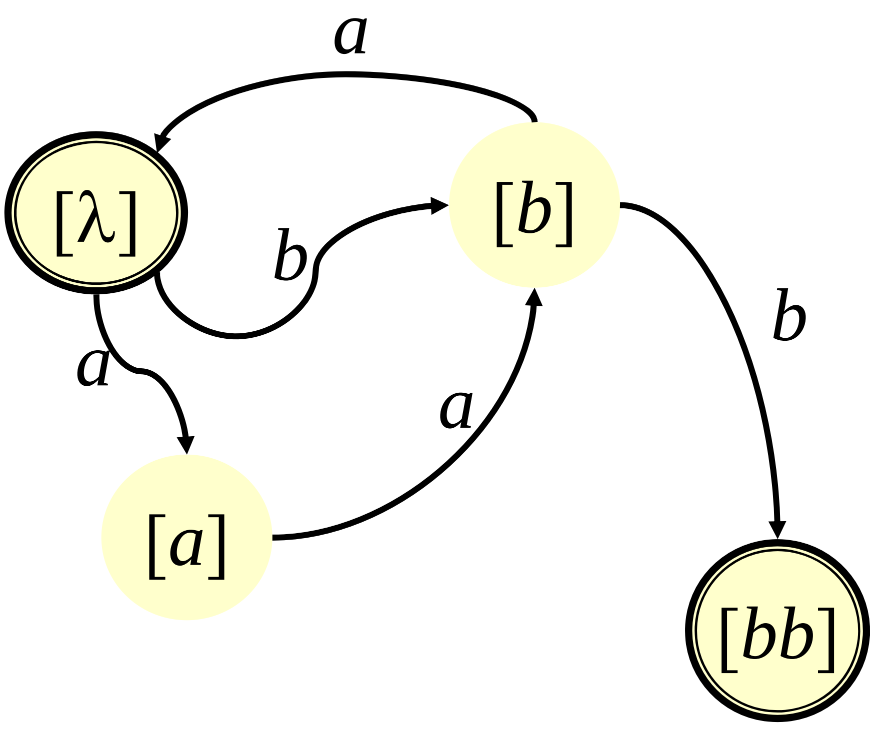
3 Gold's algorithm

E. M. Gold. Complexity of automaton identification from given data. *Information and Control*, 37:302–320, 1978.

Key ideas

- Use an observation table
- Represent the states of an automaton as strings, prefixes of the strings in the learning set
- Find some incompatibilities between these prefixes due to separating suffixes
- This leads to equivalent prefixes
- Invent the other equivalences

Strings as states



Incompatible prefixes

- $S_+=\{aab\}$
- $S_-=\{bab\}$
- Then clearly there are at least 2 states, one corresponding to *a* and another to *b*.

Observation table

- The information is organised in a table $\langle STA, EXP, OT \rangle$ where:
  - $RED \subseteq \Sigma^*$ is the *set* of states
  - $BLUE \subseteq \Sigma^*$ is the *set* of transitions

    $BLUE = (RED.\Sigma) \setminus RED$
  - $EXP \subseteq \Sigma^*$ is the *experiment set*
  - OT: $(STA = RED \cup BLUE) \times EXP \to \{0,1,*\}$ such that:

    $OT[u][e] =$ 1 if $ue \in S_+$

    0 if $ue \in S_-$

    * otherwise

An observation table

The experiments (*EXP*)

	λ	*a*
λ	1	0
a	0	*
b	1	0
aa	*	0
ab	*	*

The states (*RED*)

The transitions (*BLUE*)

Meaning

	λ	a
λ	1	0
a	0	*
b	1	*
aa	*	0
ab	*	0

$$\delta(q_0, \lambda.\lambda) \in F$$

$$\Leftrightarrow$$

$$\lambda \in L$$

	λ	a
λ	1	0
a	0	*
b	1	*
aa	*	0
ab	1	0

$$\delta(q_0, ab.a) \notin F$$

$$\Leftrightarrow$$

$$aba \notin L$$

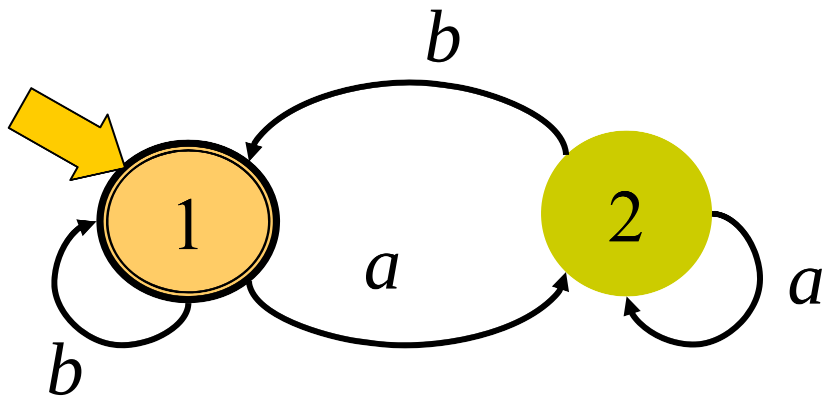
Redundancy

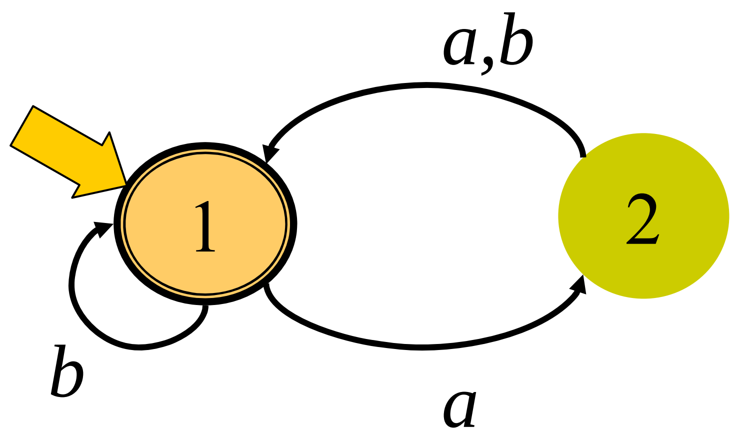
	λ	a	aa
λ	1	*	0
a	*	0	0
b	*	*	1
aa	0	0	0
ab	1	*	1

Must have identical label (redundancy)

DFA consistent with a table

	λ	*a*
λ	1	0
a	0	*
b	1	*
aa	*	0
ab	1	0





Both DFA are consistent with this table

Closed table without holes

- Let $u \in RED \cup BLUE$ and let $OT[u][e]$ denote a cell

$OT[u]$ denotes the row indexed by u

- We say that the table is *closed if*
$\forall t \in BLUE, \exists s \in RED : OT[t] = OT[s]$

- We say that the table has *no holes if*
$\forall u \in RED \cup BLUE, \forall e \in E \; OT[u][e] \in \{0,1\}$

This table is closed

	λ	a
λ	1	0
a	0	0
b	1	0
aa	0	0
ab	1	0

This table is not closed

	λ	a
λ	1	0
a	0	0
b	1	0
aa	0	1
ab	1	0

Not closed

An equivalence relation

Let $<RED,EXP,OT>$ be a closed table and with no holes. Consider the equivalence relation over STA:

$$s_1 \equiv s_2$$

if

$$OT[s_1]=OT[s_2] \wedge \forall a \in \Sigma \; OT[s_1a]=OT[s_2a]$$

Class of s is denoted by ***[s]*** *(also!)*

Equivalent prefixes

	λ	a
λ	1	0
a	0	0
b	1	0
aa	0	0
ab	1	0

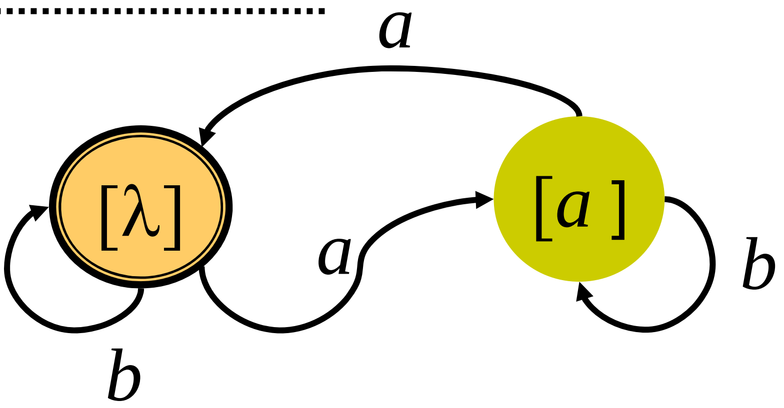
prefixes λ and b are equivalent since $OT[\lambda]=OT[b]$

Building an automaton from a table

We define $A(<STA, EXP, OT>) = (\Sigma, Q, \delta, q_0, F)$ as follows:

- $Q = \{[s]: s \in RED\}$
- $q_0 = [\lambda]$
- $F = \{[ue]: OT[u][e] = 1\}$
- $\delta([s_1], a) =$
  - $[s_2a]$ if $\exists s_2 \in [s_1]: s_2a \in RED$
  - any $[s]: s \in RED \wedge OT[s] = OT[s_1a]$

	λ	a
λ	1	0
a	0	1
b	1	0
aa	1	0
ab	0	1

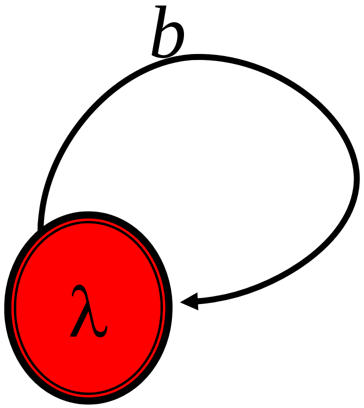


Compatibility Theorem:

- Let *<STA,EXP,OT>* be an observation table closed and with no holes
- If *STA* is *prefix-complete* and *EXP* is *suffix-complete* then *A*(*<STA,EXP,OT>*) is compatible with the data in *<STA,EXP,OT>*

Example

- $RED = [\lambda]$
- $Q = \{[\lambda]\}$
- $q_0 = [\lambda]$
- $F = \{[\lambda]\}$



	λ	b
λ	1	1
b	1	1

Exercise

- *Build a DFA from either of these two tables*

	λ	b	bbb
λ	1	1	0
b	1	1	0

	λ
λ	0
a	1
aaa	0
$aaaa$	0

Building the initial table from a sample

- Given a sample *S* and a set of strings (*RED*) prefix-complete, it is always possible to select a set of experiments *E* such that the table *<STA,E,OT>* contains **all** the information in *S*
- But usually this table is going to have holes

Obviously different rows

Let $s_1, s_2 \in RED \cup BLUE$ we say that $OT[s_1]$ is *obviously different* from $OT[s_2]$ *if*

$$\exists e \in E: OT[s_1][e], OT[s_2][e] \in \{0,1\} \text{ and } OT[s_1][e] \neq OT[s_2][e]$$

	λ	a
λ	1	0
a	0	0
b	1	0
aa	0	0
ab	0	1

If $\exists t \in BLUE$ such that $OT[t]$ is *obviously different* from any $OT[s]$, ($s \in RED$) then no filling of holes in $\langle RED, EXP, OT \rangle$ can produce a closed table.

	λ	a
λ	1	0
a	0	0
b	1	*
aa	0	*
ab	*	1

ab is OD with each *s*

Algorithm

$RED \leftarrow \{\lambda\}$

build <*RED,E,OT*> with *E* suffix-complete

while $\exists x \in BLUE$: *OT[x]* is OD **do**

    add *x* to *RED*

    update *BLUE*

    update <*STA,E,OT*>

$Q \leftarrow RED$

$q_0 \leftarrow \lambda$

$F \leftarrow \{t \in RED : OT[t][\lambda] = 1\}$

$\delta(t,a) \leftarrow$ ta if $ta \in RED$

t' if $t' \in RED$ and not OD

There can be several such *t'*

if <*STA,EXP,OT*> is incompatible with *S*, return the PTA

Example run

- $S_+=\{bb, abb, bba, bbb\}$
- $S_-=\{a, b, aa, ba, bab\}$

	λ	*a*	*b*	*aa*	*ab*	*ba*	*bb*	*abb*	*bab*	*bba*	*bbb*
λ	*	0	0	0	*	0	1	1	0	1	1
a	0	0	*	*	*	*	1	*	*	*	*
b	0	0	1	*	0	1	1	*	*	*	*

λ and *b* are OD

	λ	*a*	*b*	*aa*	*ab*	*ba*	*bb*	*abb*	*bab*	*bba*	*bbb*
λ	*	0	0	0	*	0	1	1	0	1	1
b	0	0	1	*	0	1	1	*	*	*	*
a	0	0	*	*	*	*	1	*	*	*	*
ba	0	*	0	*	*	*	*	*	*	*	*
bb	1	1	1	*	*	*	*	*	*	*	*

1) We promote line *b*

2) We expand the table, adding rows *ba* and *bb*

3) *bb* is OD

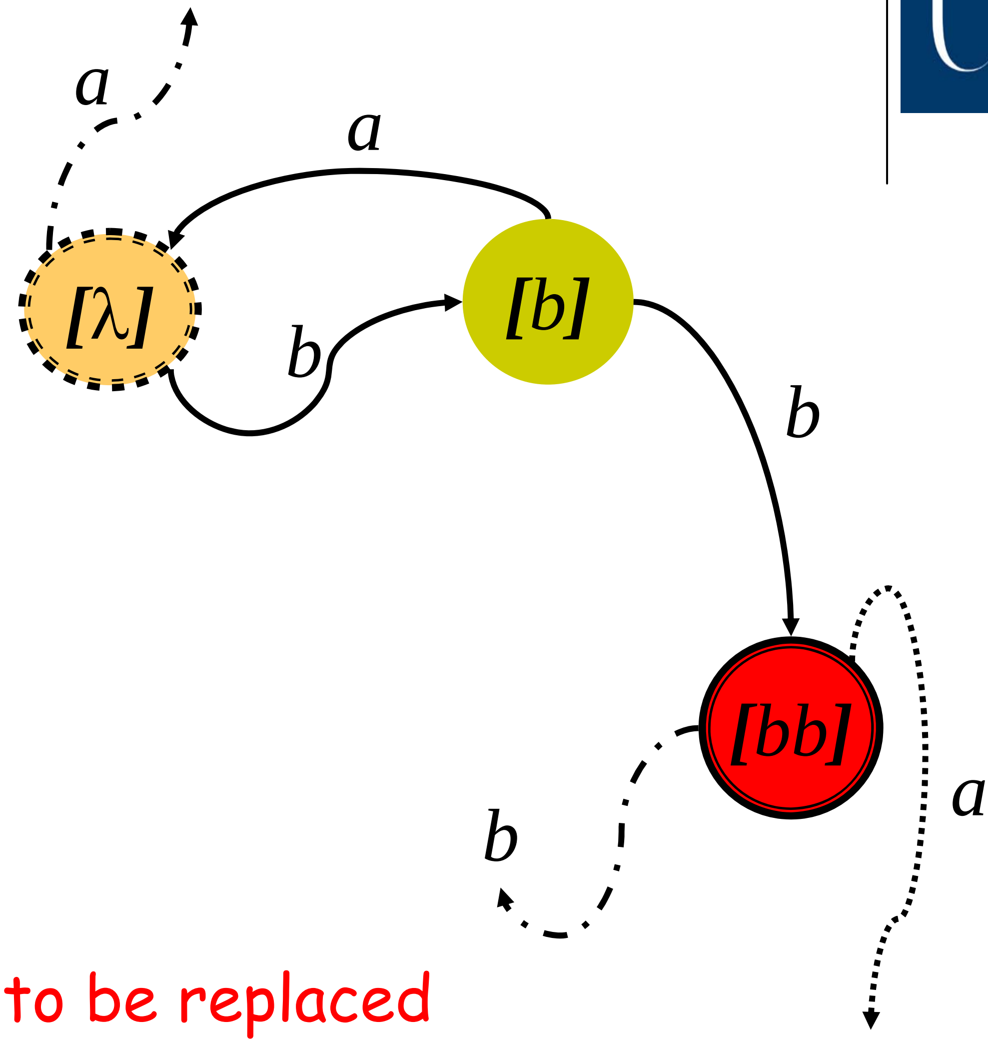
	λ	*a*	*b*	*aa*	*ab*	*ba*	*bb*	*abb*	*bab*	*bba*	*bbb*
λ	*	0	0	0	*	0	1	1	0	1	1
b	0	0	1	*	0	1	1	*	*	*	*
bb	1	1	1	*	*	*	*	*	*	*	*
a	0	0	*	*	*	*	1	*	*	*	*
ba	0	*	0	*	*	*	*	*	*	*	*
bba	1	*	*	*	*	*	*	*	*	*	*
bbb	1	*	*	*	*	*	*	*	*	*	*

1) We promote line *bb*

2) We expand the table, adding rows *bba* and *bbb*

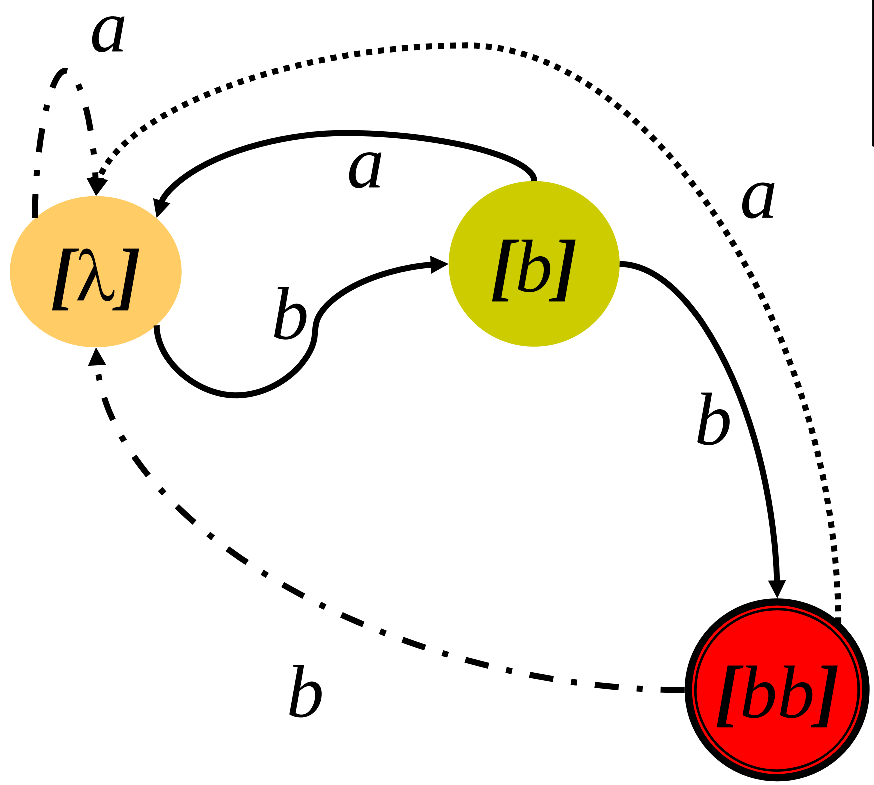
3) We construct the automaton as no line is OD

	λ	a	b
λ	*	0	0
b	0	0	1
bb	1	1	1
a	0	0	*
ba	0	*	0
bba	1	*	*
bbb	1	*	*

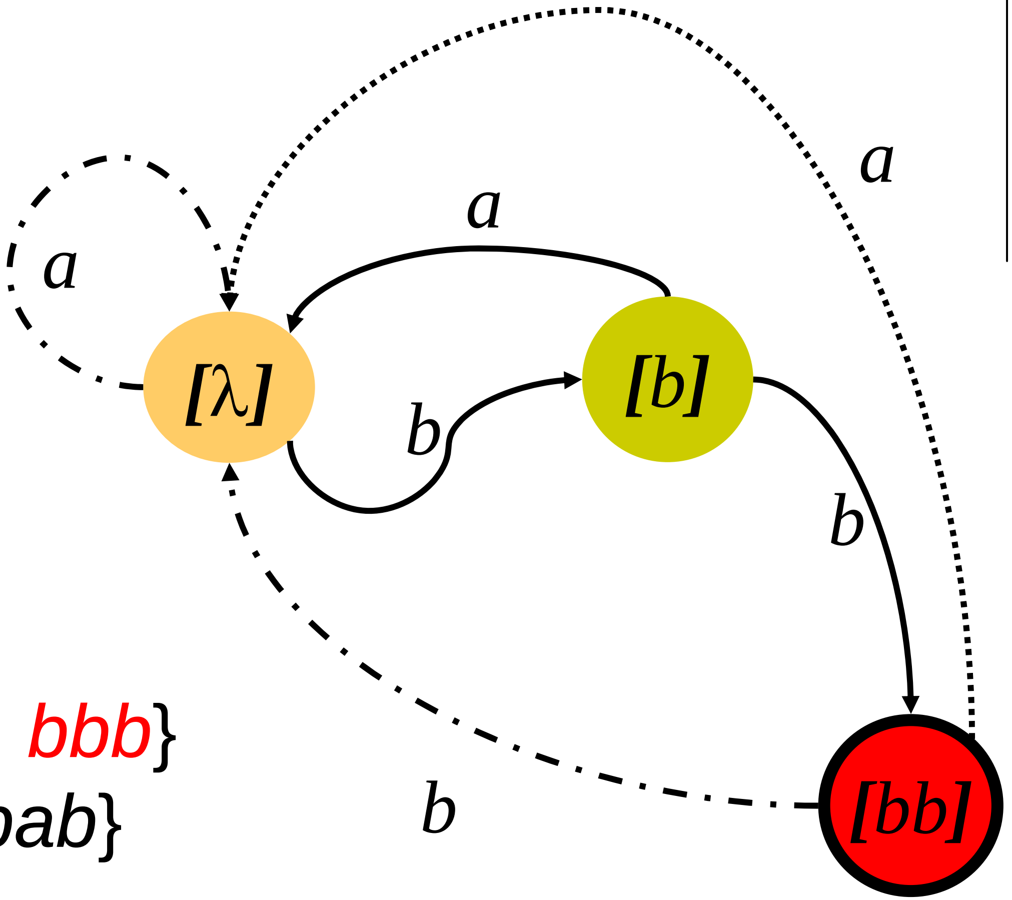


Now the * have to be replaced

	λ	a	b
λ	*	0	0
b	0	0	1
bb	1	1	1
a	0	0	*
ba	0	*	0
bba	1	*	*
bbb	1	*	*



Wild guess!



S_+={*bb*, *abb*, *bba*, *bbb*}
S_-={*a*, *b*, *aa*, *ba*, *bab*}

The automaton is inconsistent.
We shall have to return the *PTA* instead.

But !

- **a**$_{Gold}$ is deterministic: it takes deterministic decisions in order to solve the « filling holes » question
- In practice it will very often return the PTA

Equivalence of problems

Let *RED* be a state test set prefix-complete, and *S* be a sample. Let *<STA,EXP,OT>* be an observation table consistent with all the data in *S*, with *EXP* suffix-complete

The question:

Does there exist a *DFA* with the states of *RED* and compatible with *S*?

is equivalent to:

Can we fill the holes such that *<STA,EXP,OT>* is closed?

Complexity

The problem:

is there a *DFA* with states in *RED* and compatible with *S*?

is NP-Complete

Corollary

Given *S* and a positive integer *n*, the question:

Is there a *DFA* with *n* states compatible with *S*?

is NP-Complete

Properties of $\mathbf{a}_{Gold}$

1) the output is consistent with a sample S
2) $\mathbf{a}_{Gold}$ identifies in the limit any regular language
3) $\mathbf{a}_{Gold}$ works in time polynomial in $|S|$
4) if the size of the target is n, then there is a characteristic sample CS with $|CS|= 2n^2(|\Sigma|+1)$, such that $\mathbf{a}_{Gold}(S)$ produces the canonical acceptor for all $S \supseteq CS$

Open questions

- Can one fill the holes in a more “intelligent” way?
- How fast can we detect that a choice (for a filling) is good or bad?

Exercise

- Run Gold's algorithm for the following data:
- $S_+=\{a, abb, bab, babbb\}$
- $S_-=\{ab, bb, aab, b, aaaa, babb\}$

4 RPNI

Regular Positive and Negative Grammatical Inference

Inferring regular languages in polynomial time. Jose Oncina & Pedro García. *Pattern recognition and image analysis, 1992*

- RPNI is a state merging algorithm
- RPNI identifies any regular language in the limit
- RPNI works in polynomial time
- RPNI admits polynomial characteristic sets

A=PTA(S+); *Blue* ={$\delta(q_I,a)$: $a\in\Sigma$ };

Red ={q_I}

While *Blue*$\neq\emptyset$ do

choose q from *Blue*

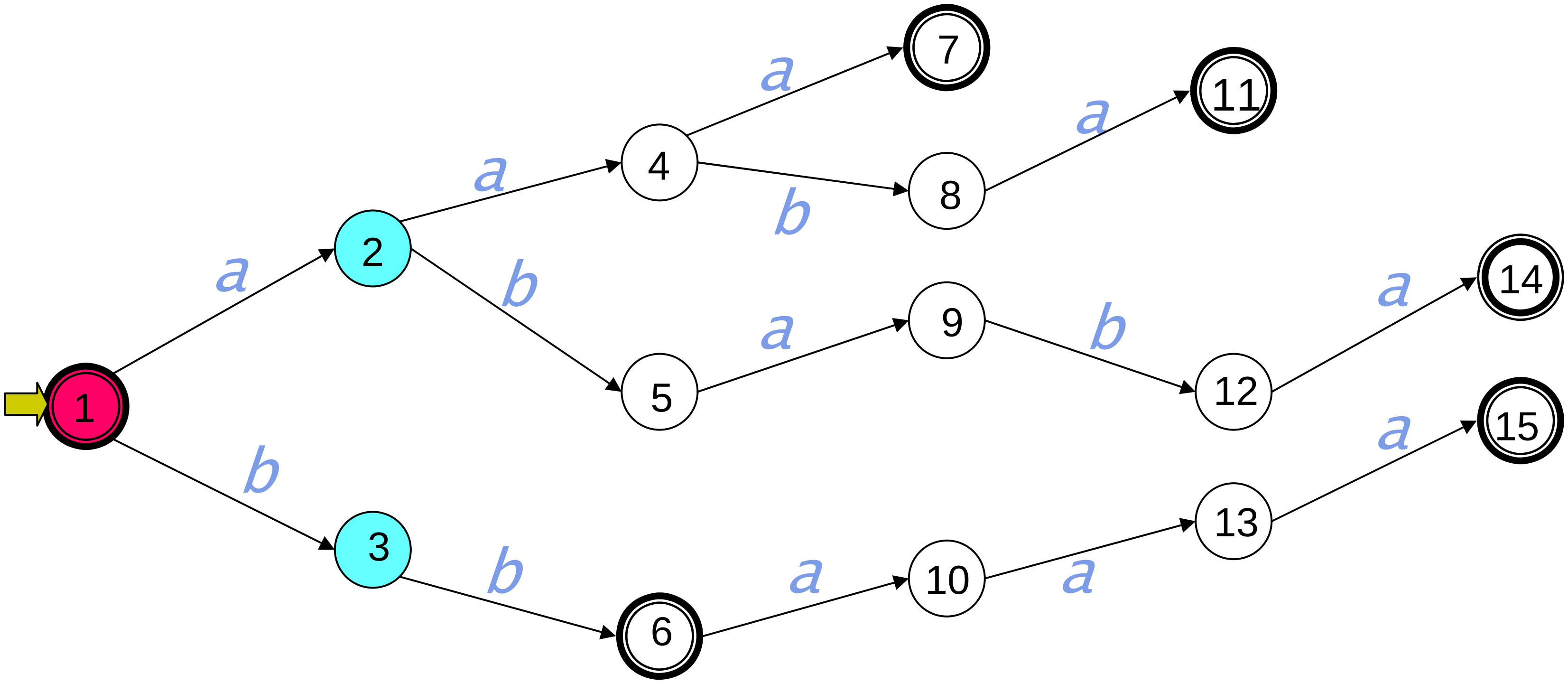
if $\exists p\in$ *Red*: $\mathbb{L}$(merge_and_fold(A,p,q))$\cap S_- = \emptyset$

then A = merge_and_fold(A,p,q)

else *Red* = *Red* $\cup$ {q}

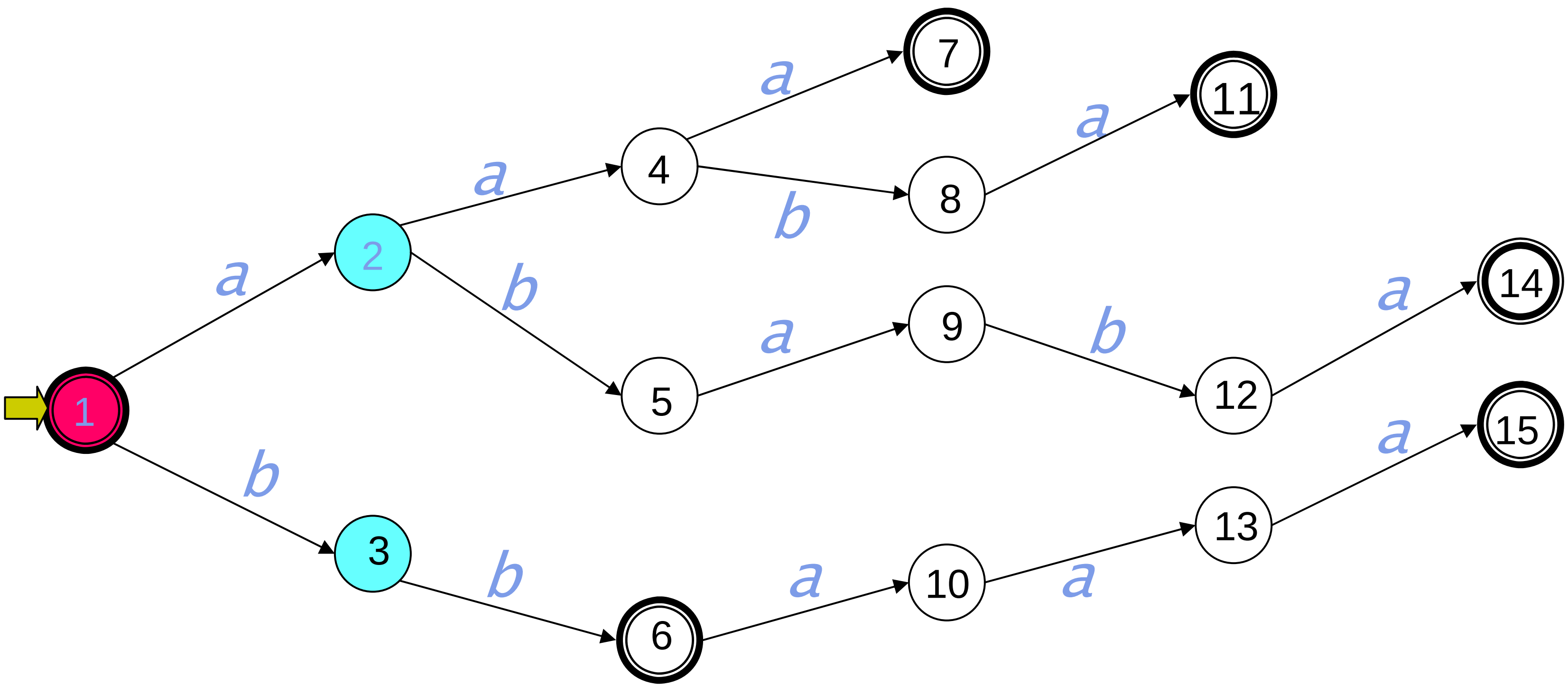
Blue = {$\delta(q,a)$: $q\in$ *Red*} - {*Red*}

$S_+ = \{\lambda,\ aaa,\ aaba,\ ababa,\ bb,\ bbaaa\}$



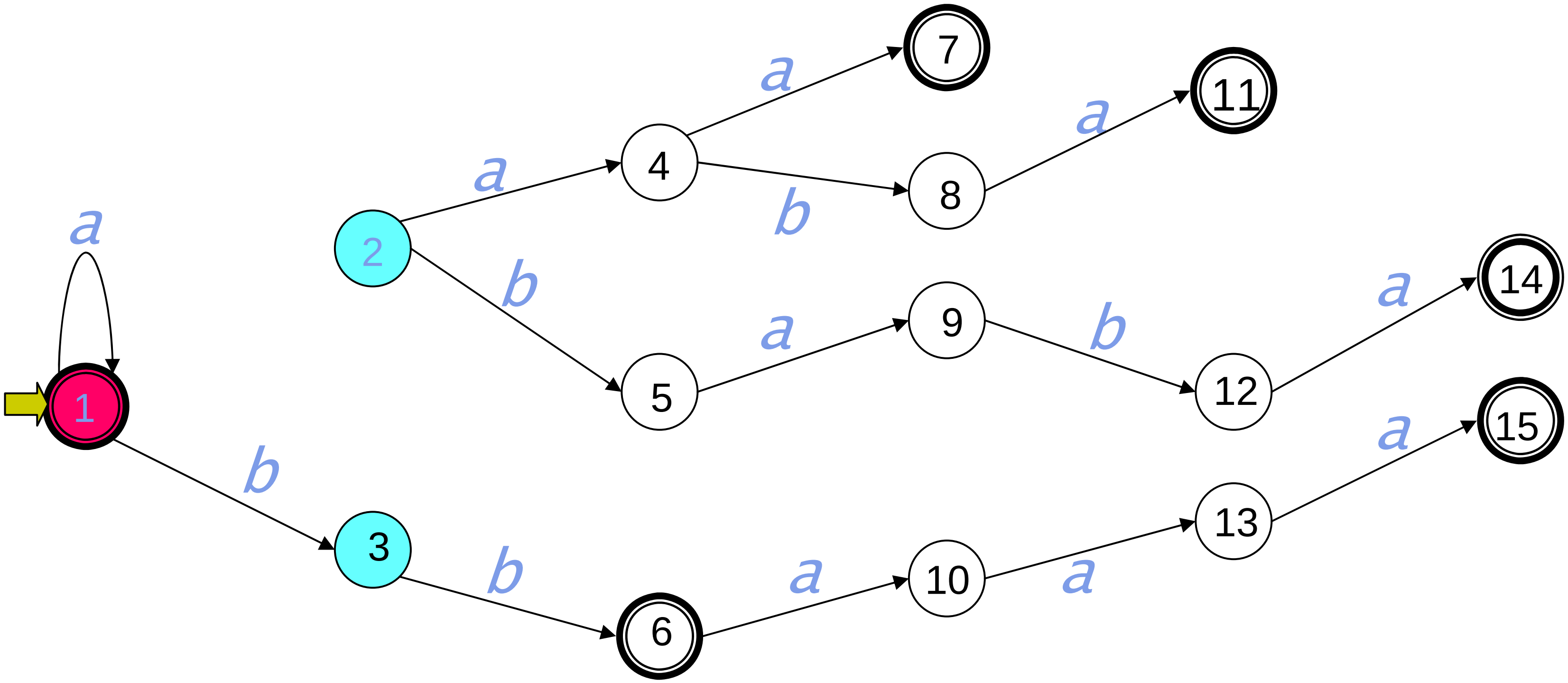
$S_- = \{aa,\ ab,\ aaaa,\ ba\}$

Try to merge 2 and 1



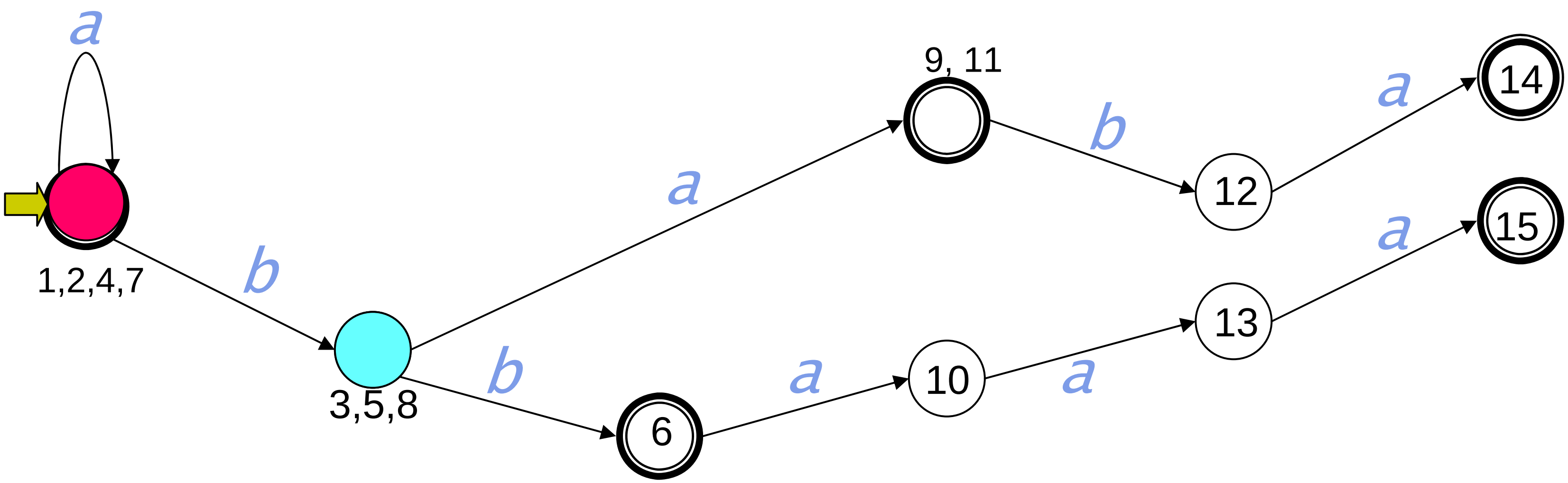
$S_{-}=\{aa, ab, aaaa, ba\}$

First merge, then fold



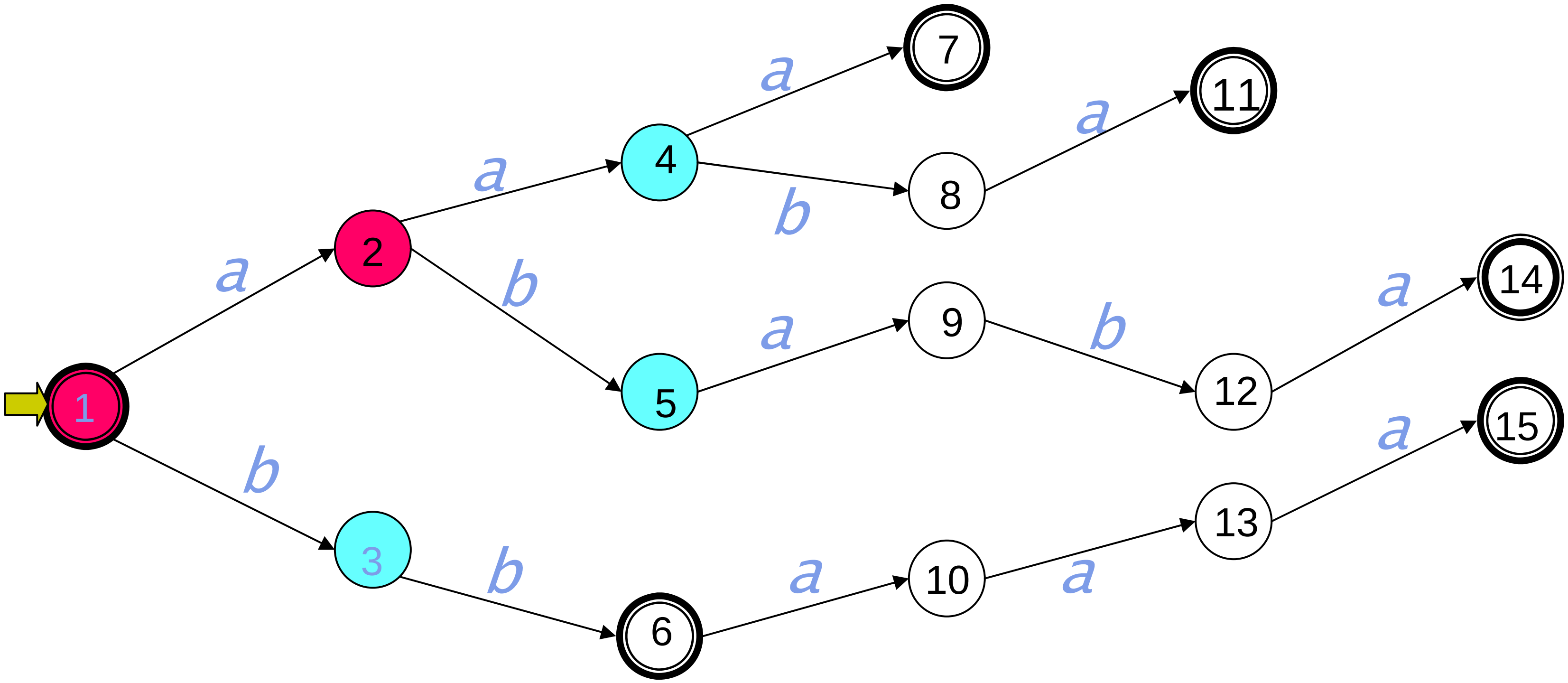
$S_{-}=\{aa,\ ab,\ aaaa,\ ba\}$

But now string aaaa is accepted, so the merge must be rejected, and state 2 is promoted



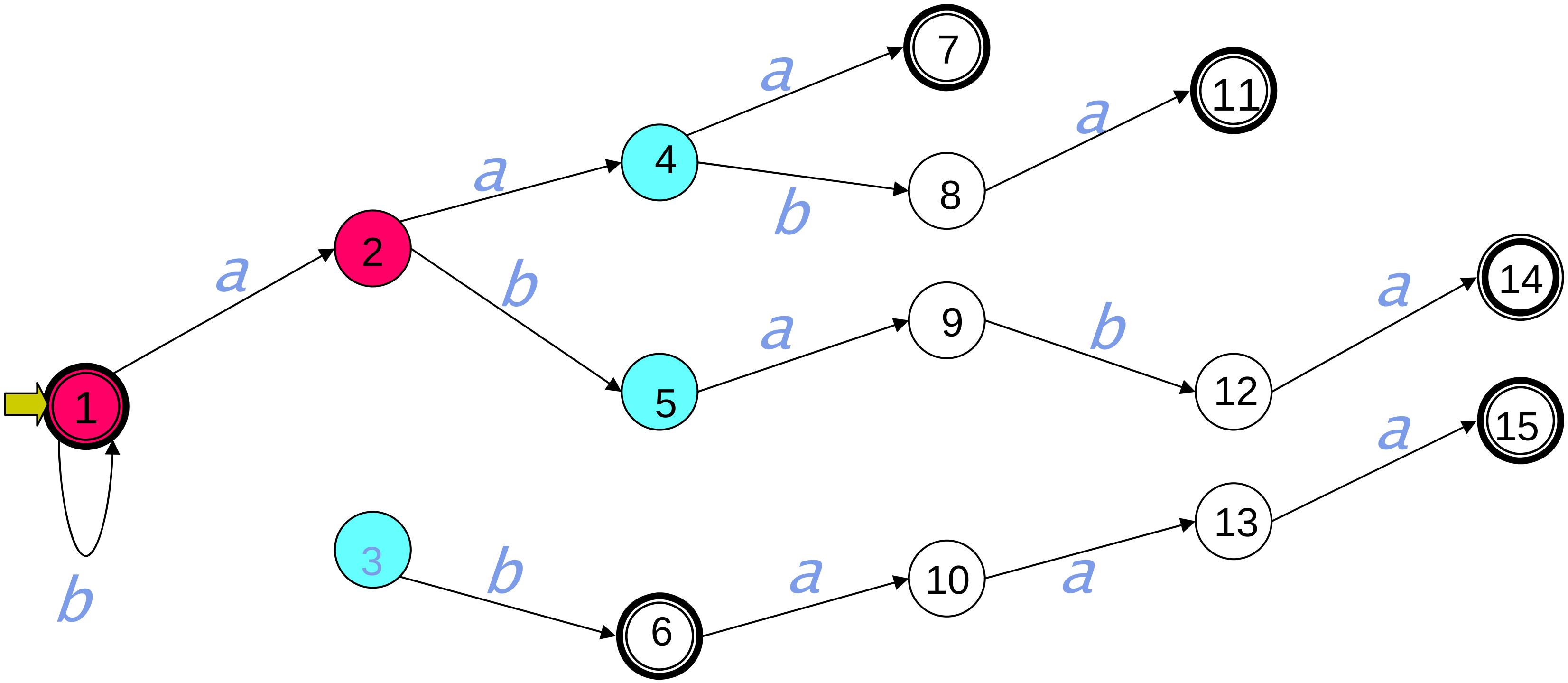
S_{-}={*aa*, *ab*, *aaaa*, *ba*}

Try to merge 3 and 1



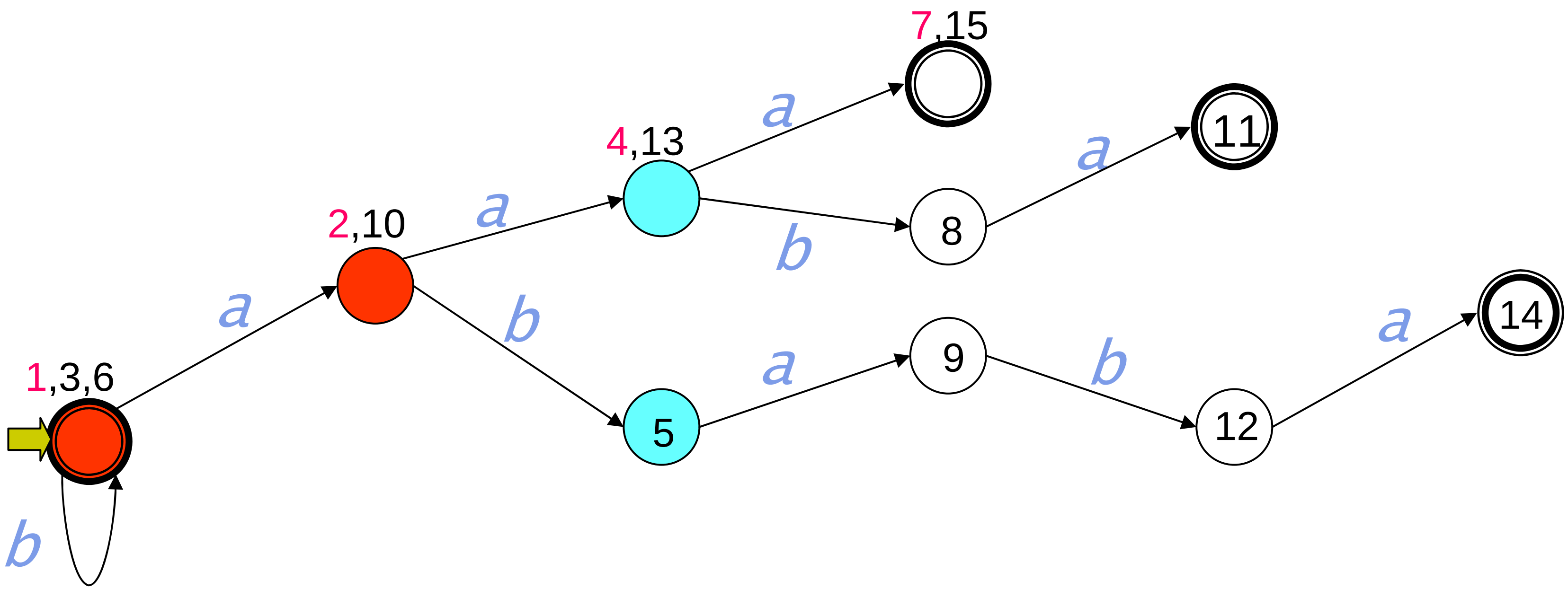
$S_- = \{aa, ab, aaaa, ba\}$

First merge, then fold



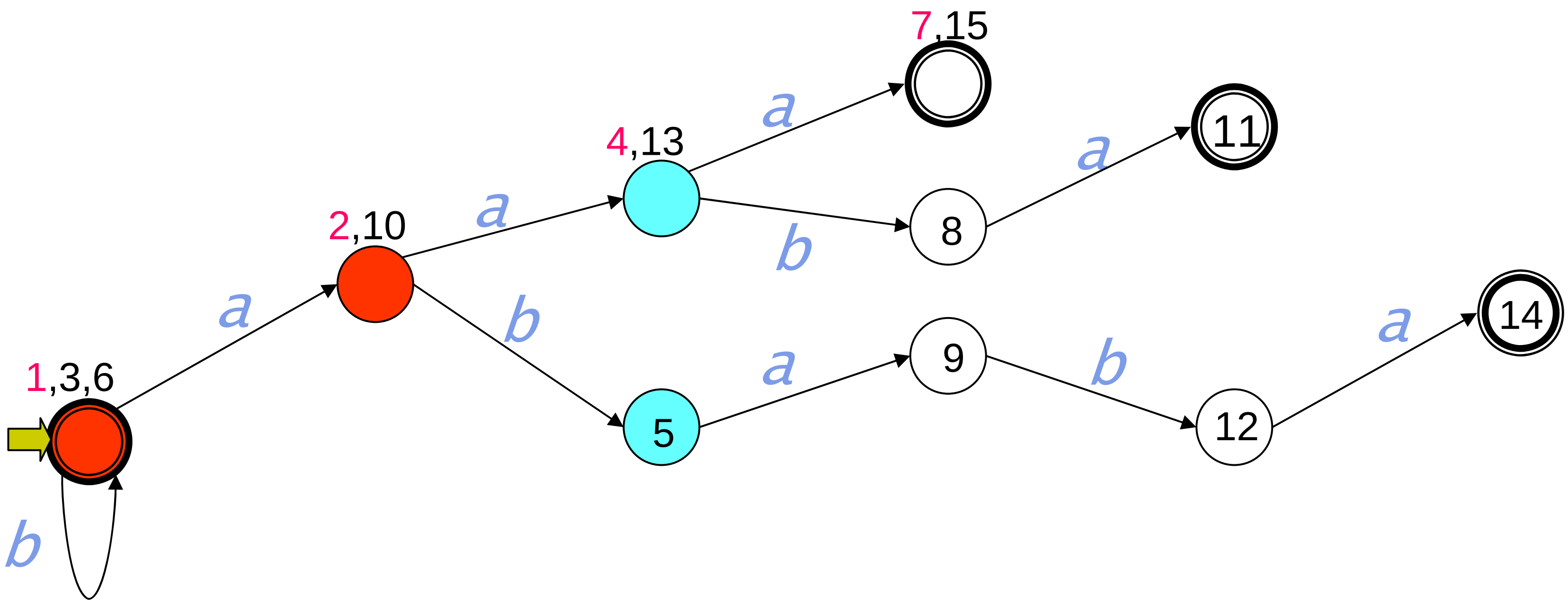
$S_{-}=\{aa,\ ab,\ aaaa,\ ba\}$

No counter example is accepted so the merge is kept



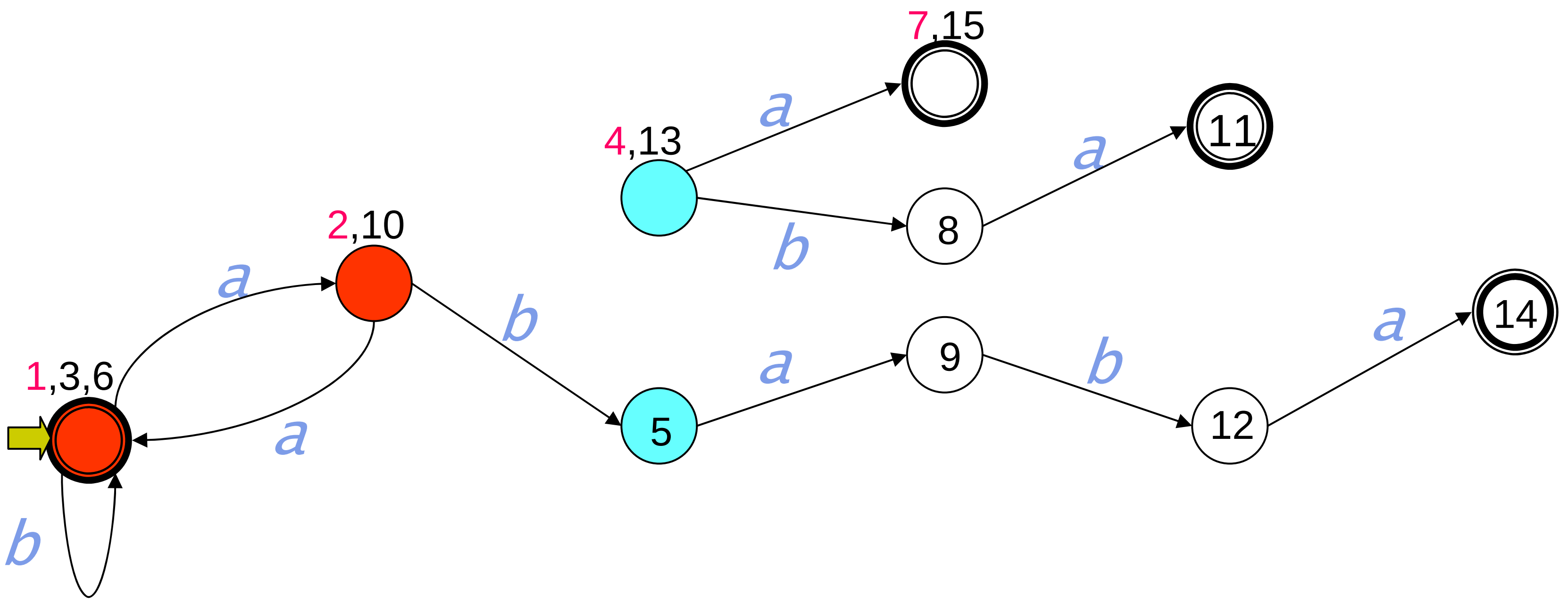
$S_- = \{aa, ab, aaaa, ba\}$

Next possible merge to be checked is {4,13} with {1,3,6}



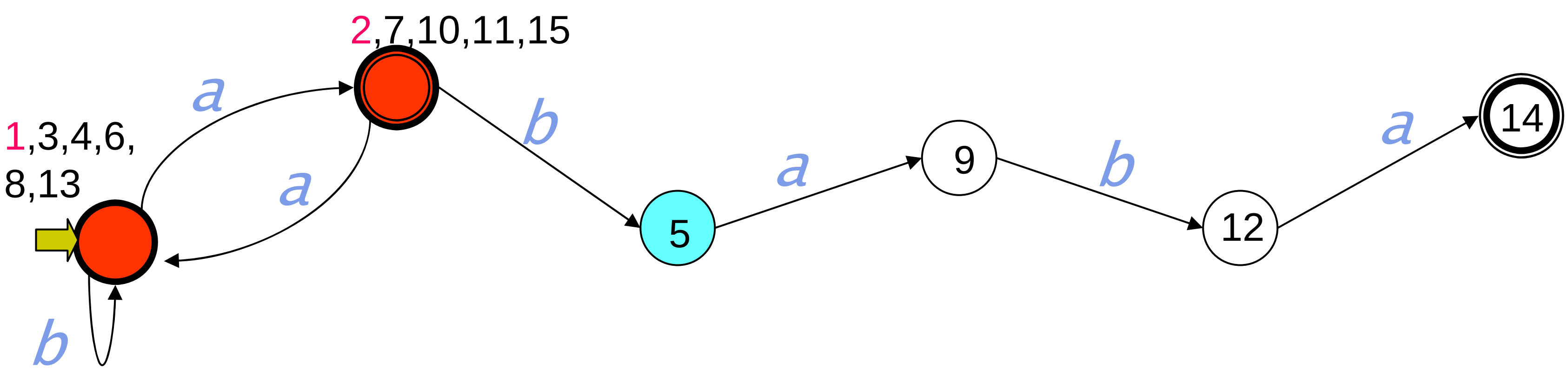
$S_- = \{aa,\ ab,\ aaaa,\ ba\}$

Merged. Needs folding subtree in {4,13} with {1,3,6}



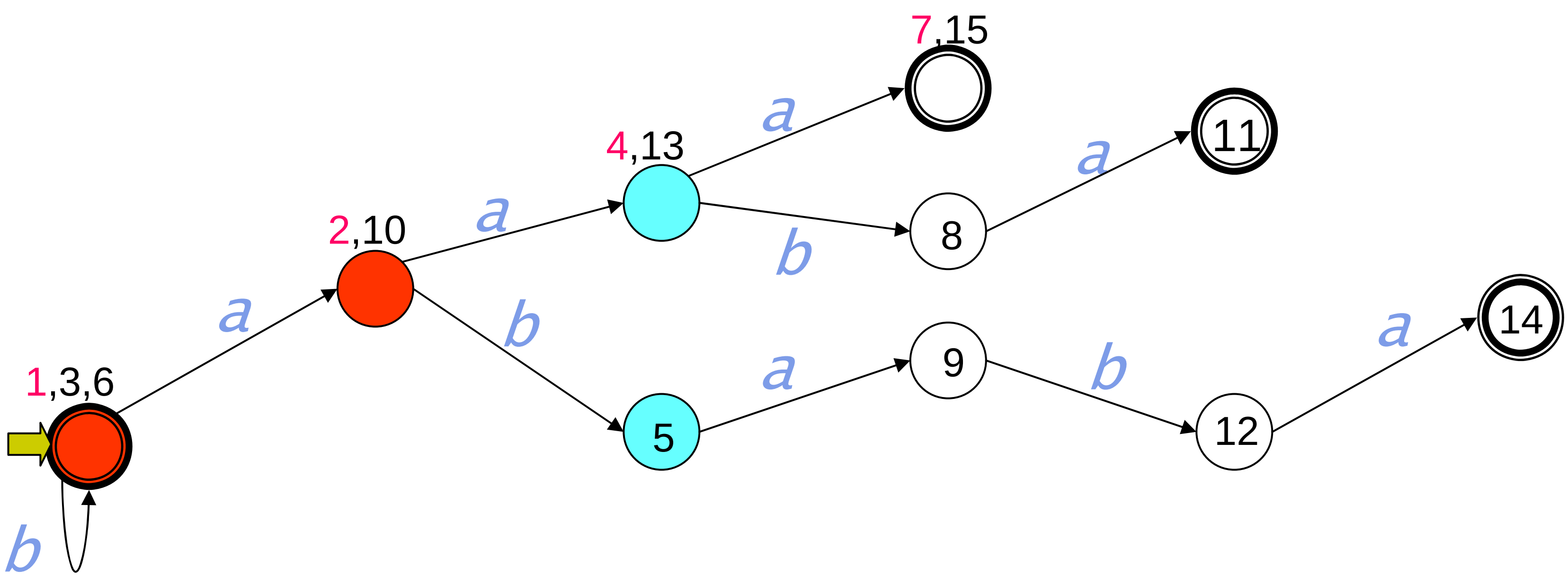
$S_- = \{aa,\ ab,\ aaaa,\ ba\}$

But now aa is accepted



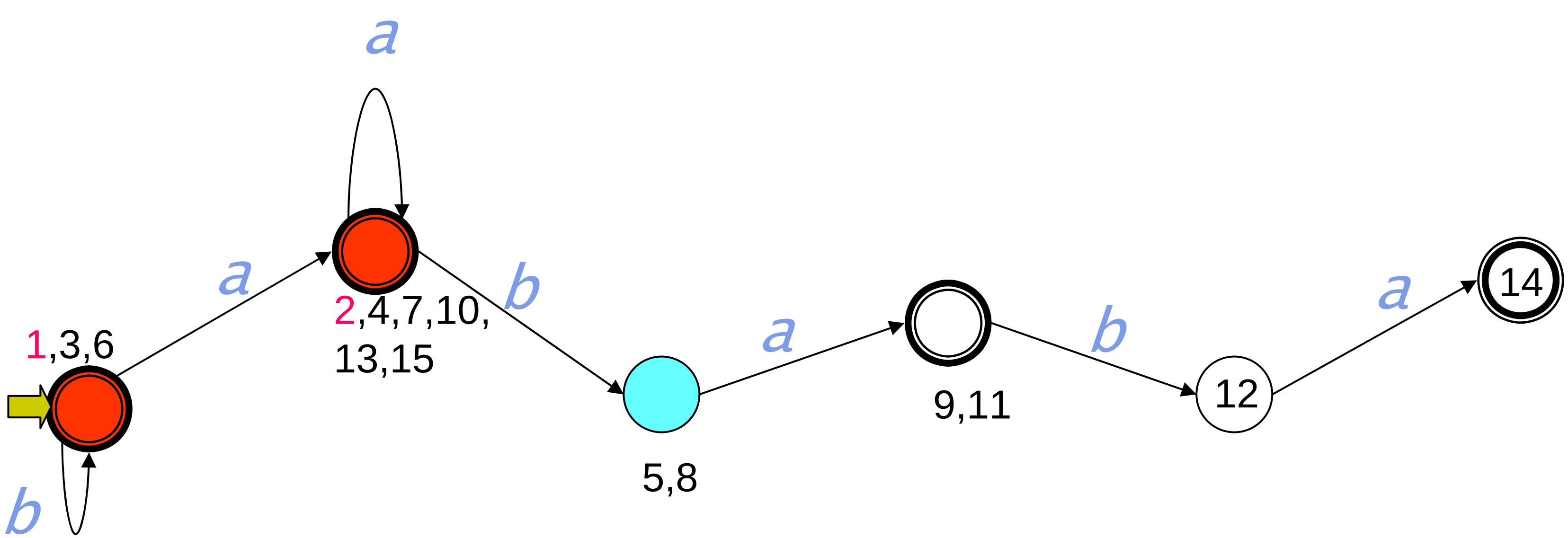
S_{-}={*aa*, *ab*, *aaaa*, *ba*}

So we try {4,13} with {2,10}



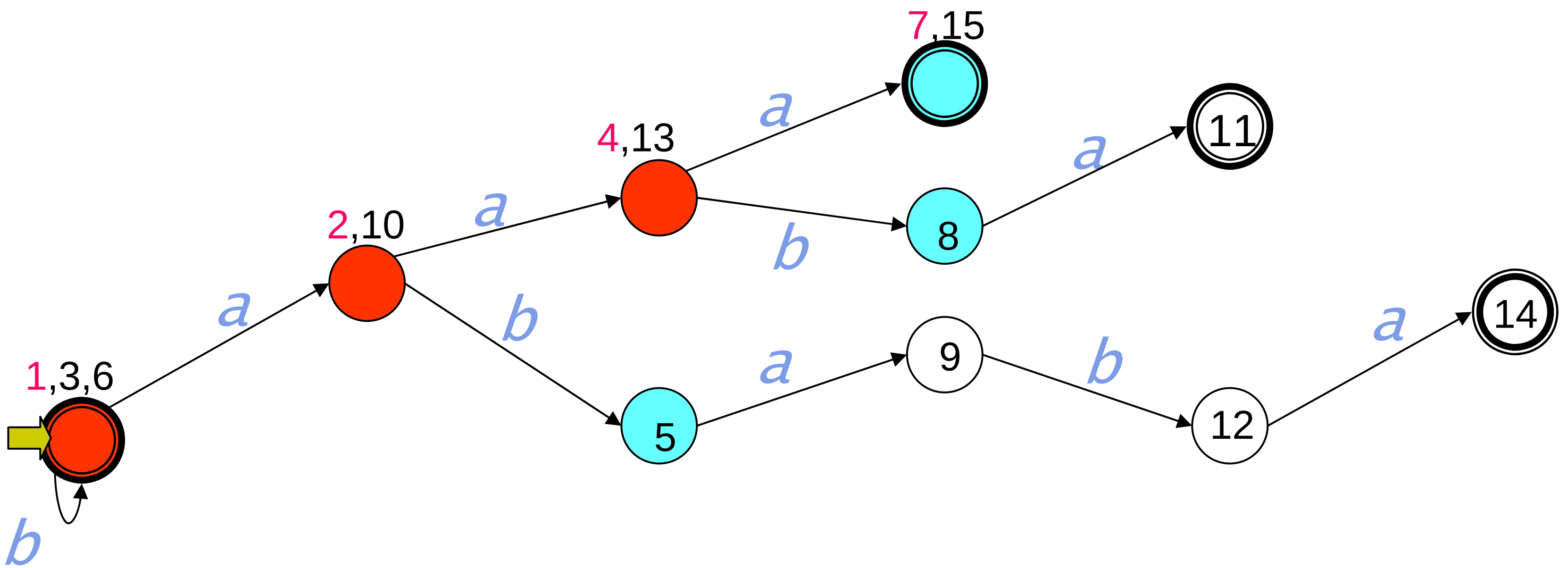
S_-={*aa*, *ab*, *aaaa*, *ba*}

Negative string aa is again accepted. Since we have tried all Red for merging, state 4 is promoted.



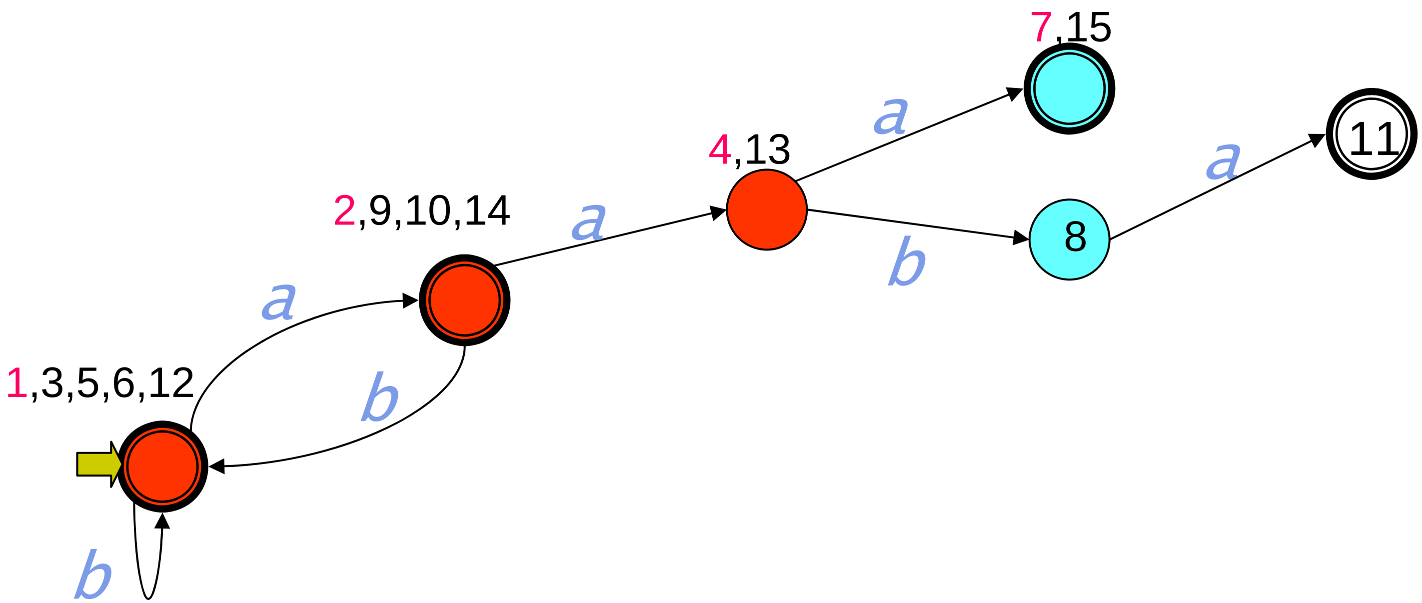
$S_- = \{aa, ab, aaaa, ba\}$

So we try 5 with {1,3,6}



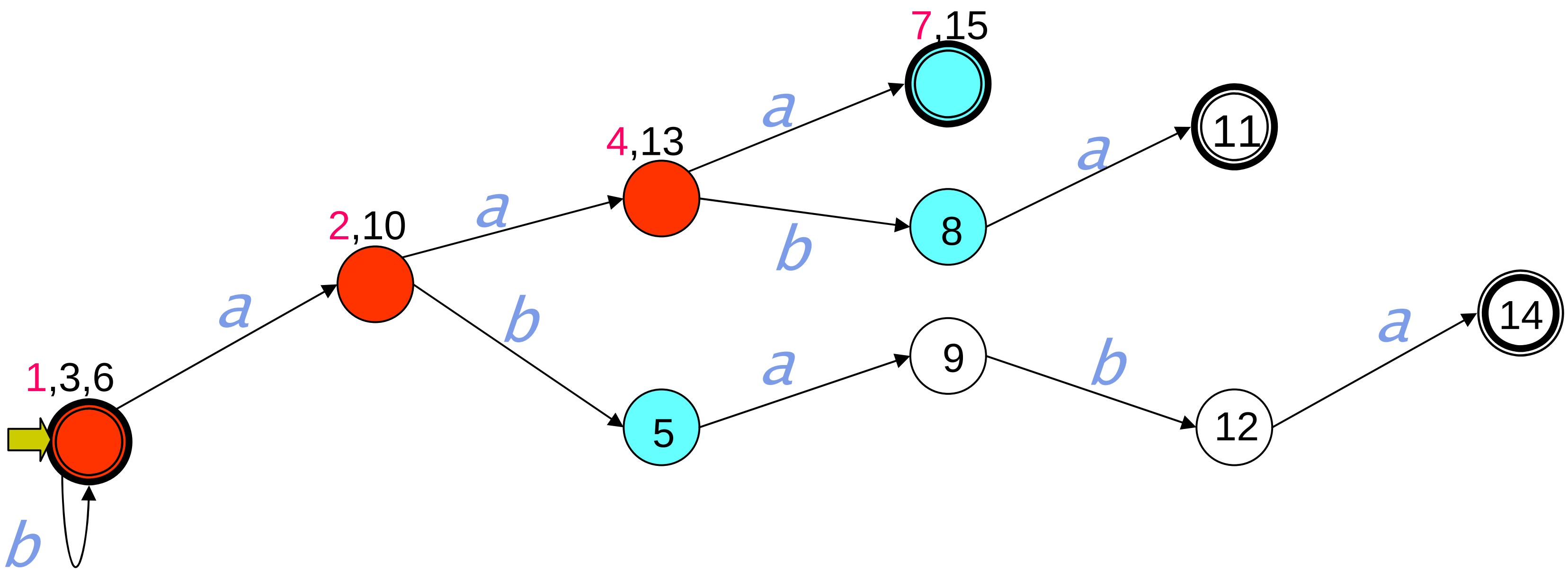
$S_- = \{aa, ab, aaaa, ba\}$

But again we accept ab



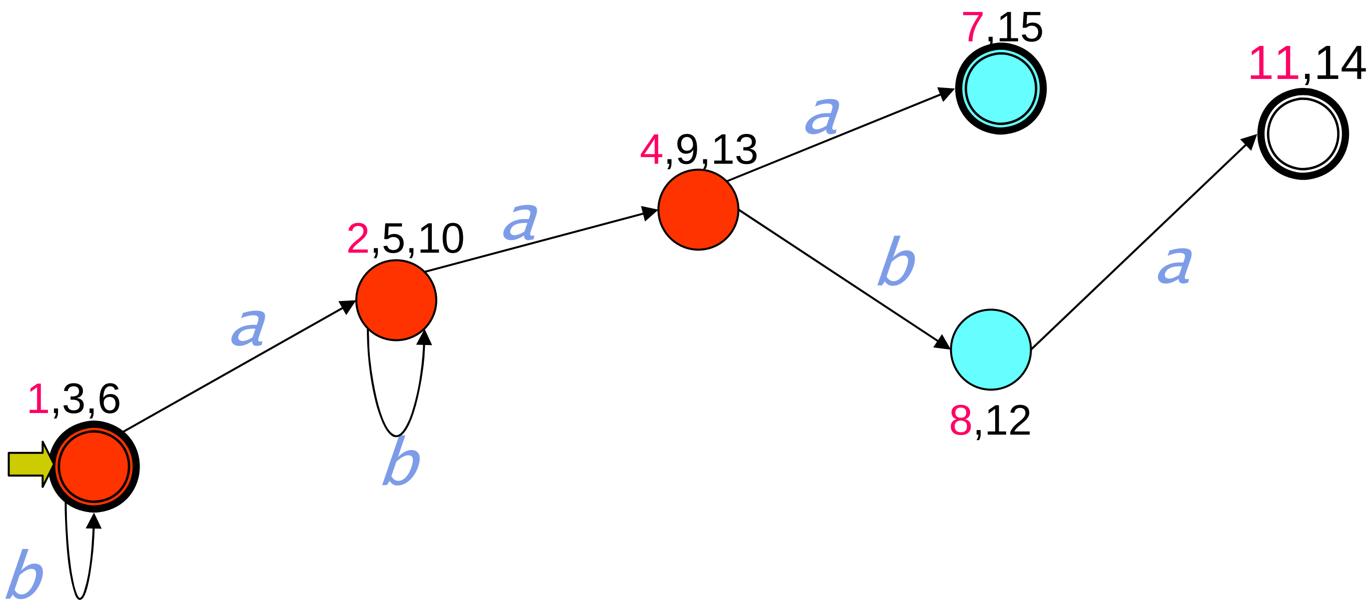
$S_- = \{aa, ab, aaaa, ba\}$

So we try 5 with {2,10}



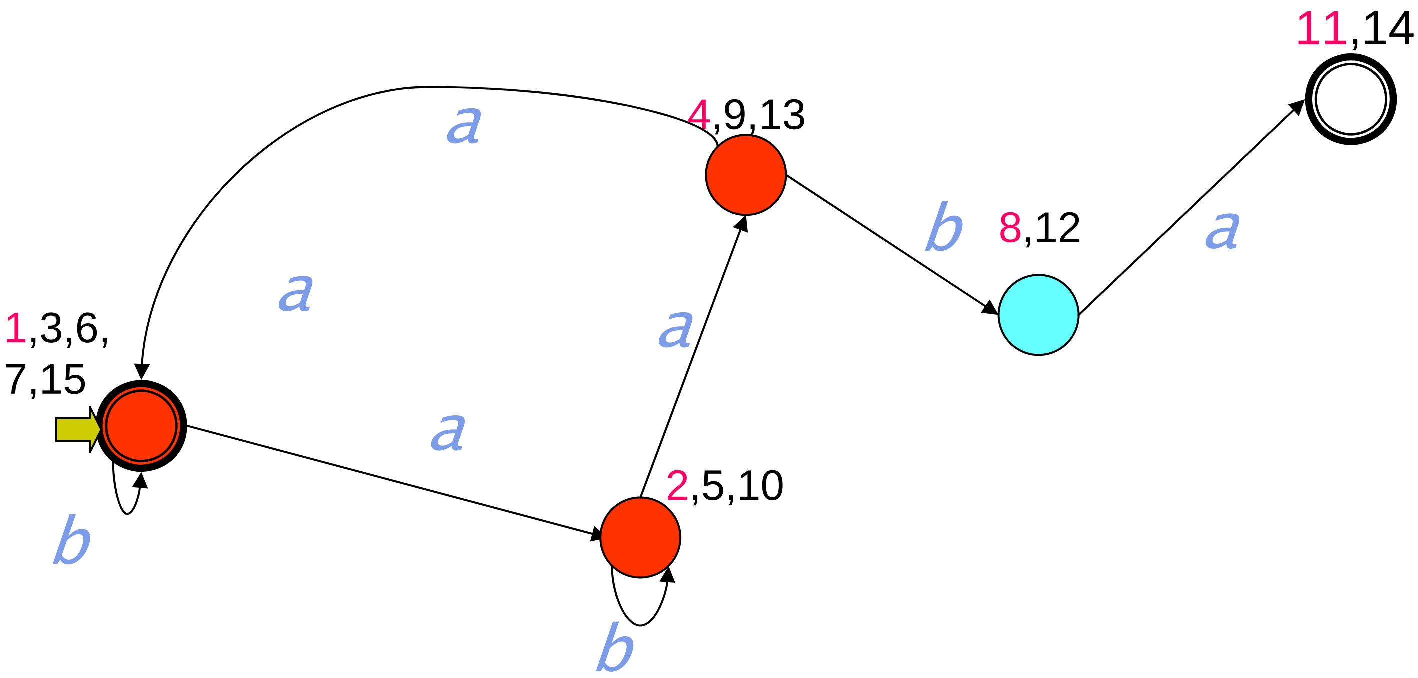
S_{-}={*aa*, *ab*, *aaaa*, *ba*}

Which is OK. So next possible merge is {7,15} with {1,3,6}



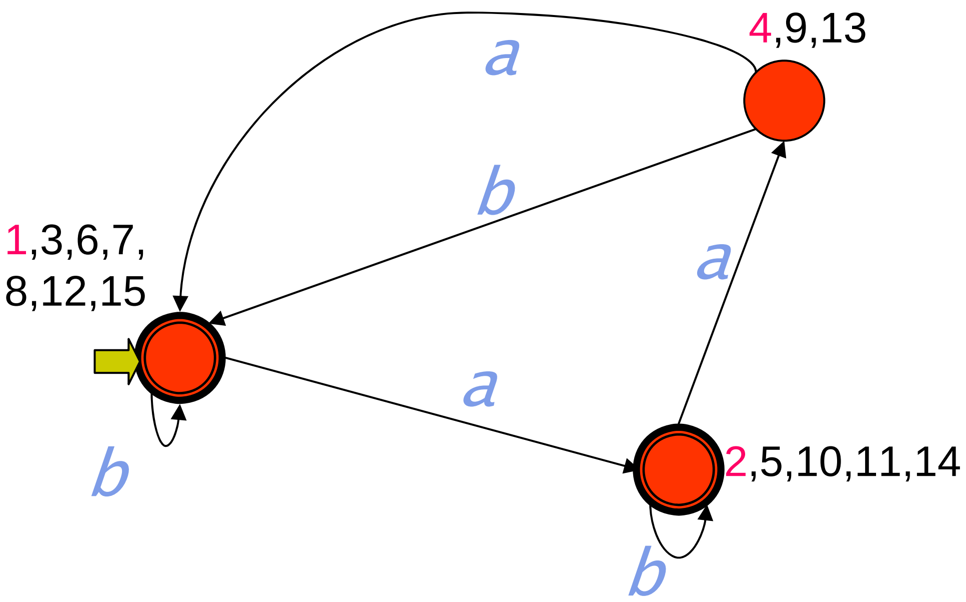
S_-={aa, ab, aaaa, ba}

Which is OK. Now try to merge {8,12} with {1,3,6,7,15}



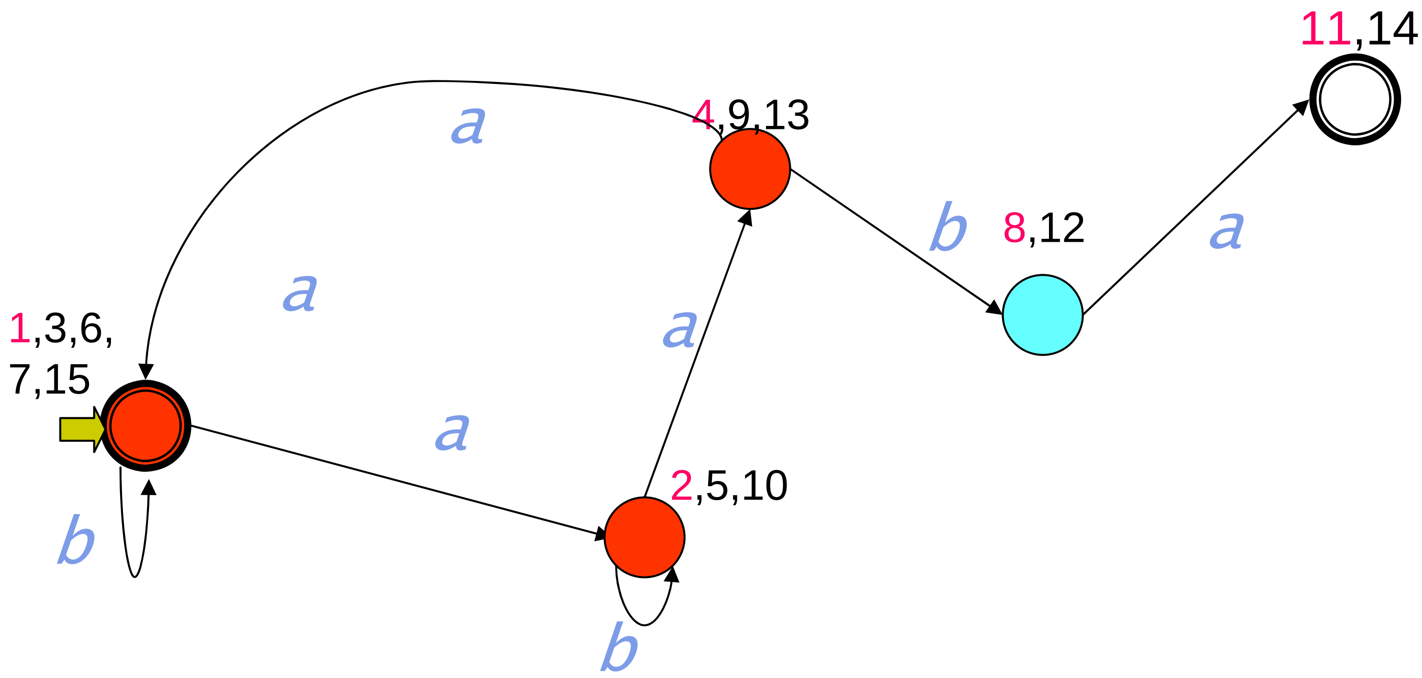
$S_- = \{aa, ab, aaaa, ba\}$

And ab is accepted



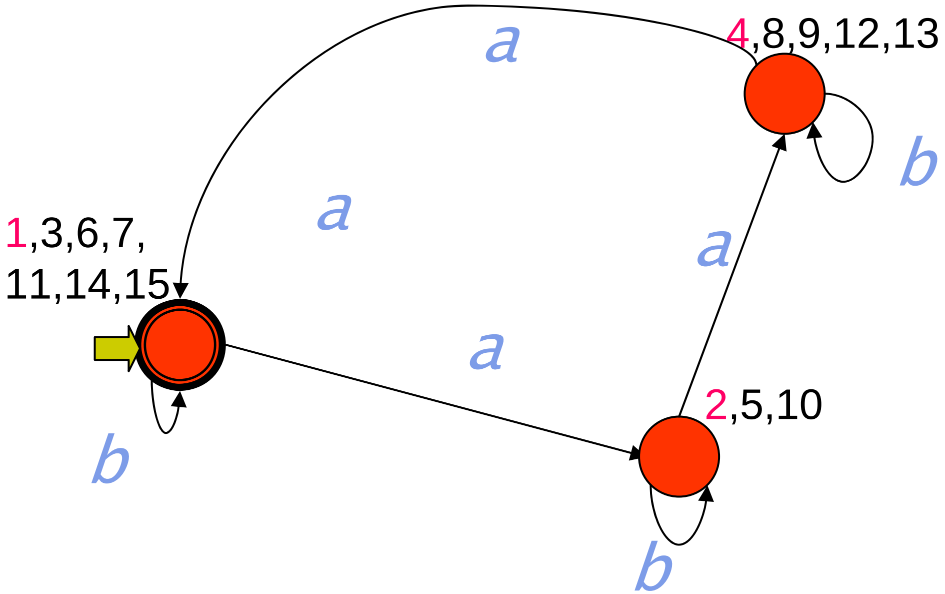
S_{-}={*aa*, *ab*, *aaaa*, *ba*}

Now try to merge {8,12} with {4,9,13}



$S_- = \{aa, ab, aaaa, ba\}$

This is OK and no more merge is possible so the algorithm halts



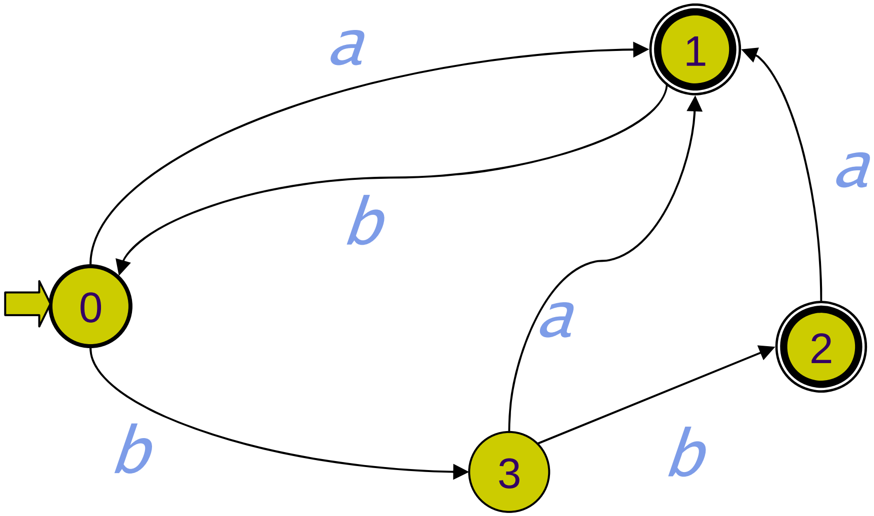
$S_- = \{aa, ab, aaaa, ba\}$

Properties

- RPNI identifies any regular language in the limit
- RPNI works in polynomial time. Complexity is in $O(\| S_+ \|^3 . \| S_- \|)$
- There are many significant variants of RPNI
- RPNI can be extended to other classes of grammars

Exercices

- Run RPNI on
  - $S_+=\{a,bba,bab,aabb\}$
  - $S_-=\{b,ab,baa,baabb\}$
- Find a characteristic sample for:



5 Complexity issues

A characteristic sample

- A sample is characteristic (for some algorithm) whenever, when included in the learning sample, the algorithm returns the correct DFA
- The characteristic sample should be of polynomial size
- There is an algorithm which given a DFA builds a characteristic sample

Definition : polynomial characteristic sample

$\mathcal{G}$ has polynomial characteristic samples for identification algorithm **a** if there exists a polynomial $p()$ such that: given any G in $\mathcal{G}$, $\exists CS$ correct sample for G, such that when $CS \subseteq f_n$, $\mathbf{a}(f_n) \equiv G$ and $\| CS \| \leq p(\| G \|)$

About characteristic samples

- If you add more strings to a characteristic sample it still is characteristic
- There can be many different characteristic samples (EDSM, tree version,…)
- Change the ordering (or the exploring function in RPNI) and the characteristic sample will change

Open problems

- RPNI's complexity is not a tight upper bound. Find the correct complexity
- The definition of the characteristic sample is not tight either. Find a better definition
- Can there be a linear time DFA learner?

Collusion

- Collusion consists in having the learner and the teacher agree of some specific encoding system. Then, the teacher can just pass one string which is the encodinng of the target.
- Is that cheating?
- Is that learning?

6 Heuristics

6.1 Genetic Algorithms

- The principle: via evolutionary mechanisms nature increases the quality of its population.
- Allow a population of solutions to interact and evolve.

Mechanisms (gene level):

- Mutation
- Crossing-over

(a solution is just a string)

Mutation

TTAGCCTTC

TTTGCCTTC

Crossing-over

- TTATCCGT TAGGCTTC

TTATC CGT TAGG CTTC

TTATC CTTC TAGG CGT

TTATCCTTC TAGGCGT

Idea: define the solutions as sequences

- Be able to measure the quality of a solution
- Conceive a first generation
- Define the genetic operations (mutation, crossing over)
- Keep the best elements of the second generation
- Iterate

Genertic algorithms in Grammatical Inference

- (Dupont 94)
  - code the automata (the partition of states of $PTA(S_+)$) into partitions
  - define genetic operators
  - define an optimum as an automaton with as few states as possible and rejecting S_-
  - run the genetic algorithm

Structural Mutation

- Select a state from a block and move it to another block
- Example:

{{1,3} {2} {4,5}}

{{1} {2} {3,4,5}}

{{1} {2} {3} {4,5}}

Structural crossover

{1,2}{3,4,5}   {1,3}{4}{2,5}

{1,3}{3,4,5}   *{1,2}{4}{2,5}*

{1,2,3}{4,5}   {1,2,3}{4}{5}

Group number encoding

Partition

{{1,2,6}{3,7,9,10}{4,8,12}{5}{11}}

is encoded by

(112341232253)

6.2 Tabu search

- (Giordano 96, based on Glover 89)
- General idea: search a space by choosing a point, and going to its best neighbor that is not in the tabu list.

$R \leftarrow$ the set of rules of the grammars in the search space

$G \leftarrow$ an initial grammar

$G^* \leftarrow G$ the best solution reached so far

$T \leftarrow \emptyset$ the Tabu list that cannot occur

$k \leftarrow 0$ the iterations counter

While $k \neq kmax$ **do**

select r in $R \setminus T$, such that the addition or deletion of r from G realizes the maximum of val on X

add or delete r from G

if $val(G) > val(G^*)$ **then** $G^* \leftarrow G$

Update T

$k \leftarrow k+1$

Return G^*

- Procedure Update(T, r)

  if *card*(T) = n **then** delete its last element

  Add r as the first element of T

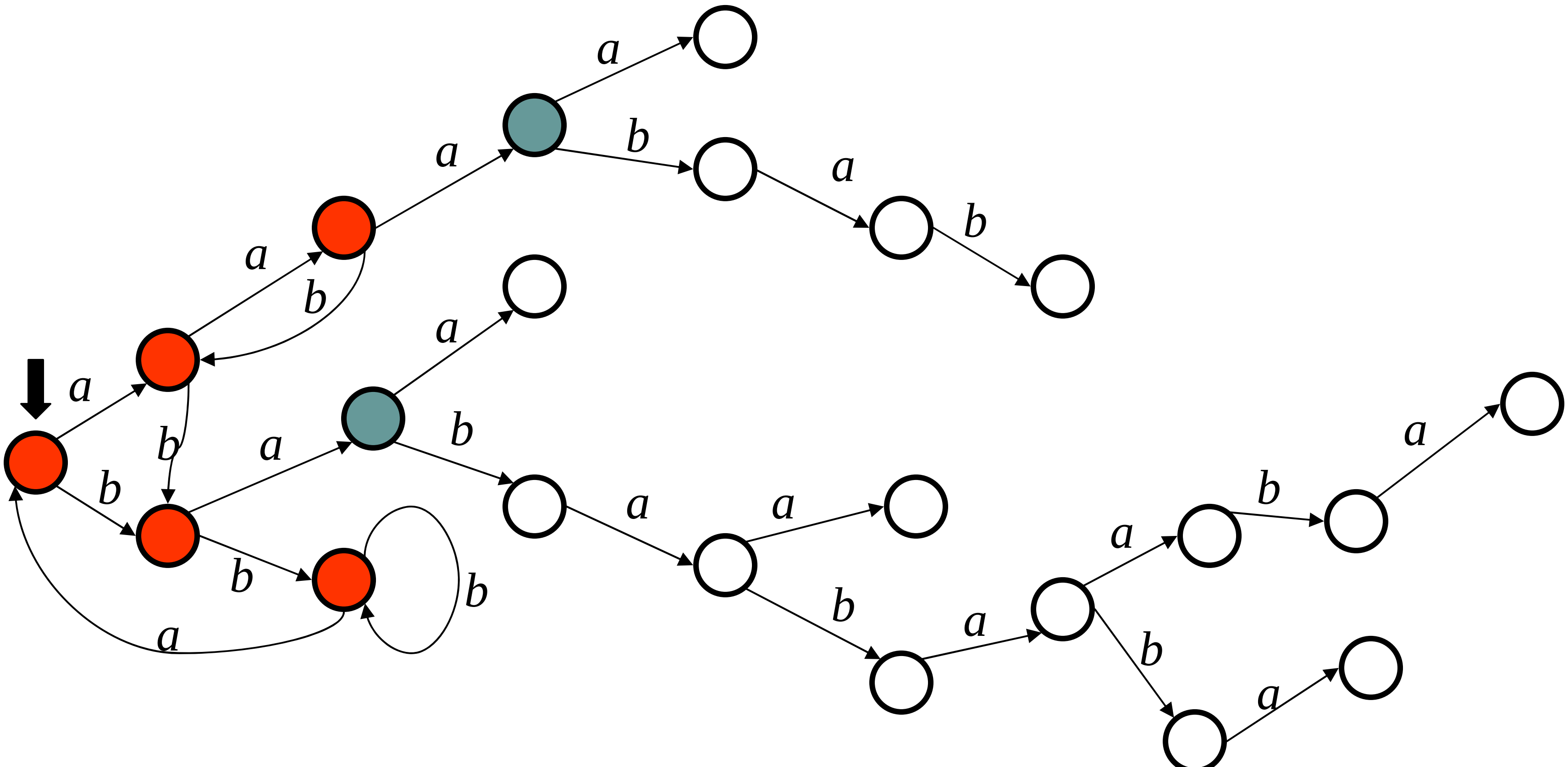
- Tricks
  - **If** *blocked* **then** delete oldest rule
  - *blocked* $\leftarrow$ 6 iterations
  - **if** new G^* **then** empty(T)

6.3 Heuristic greedy State Merging

- RPNI chooses to merge the first 2 states that can be merged
- This is an optimistic view
- There may be another…
- But remember: RPNI identifies in the limit!

How do greedy state merging algorithms work?

- choose two states
  - perform a cascade of forced merges to get a deterministic automaton
  - if it accepts sentences of $S-$, backtrack and choose another couple
  - if not, loop until no merging is still possible



The blue fringe (Lang 98)

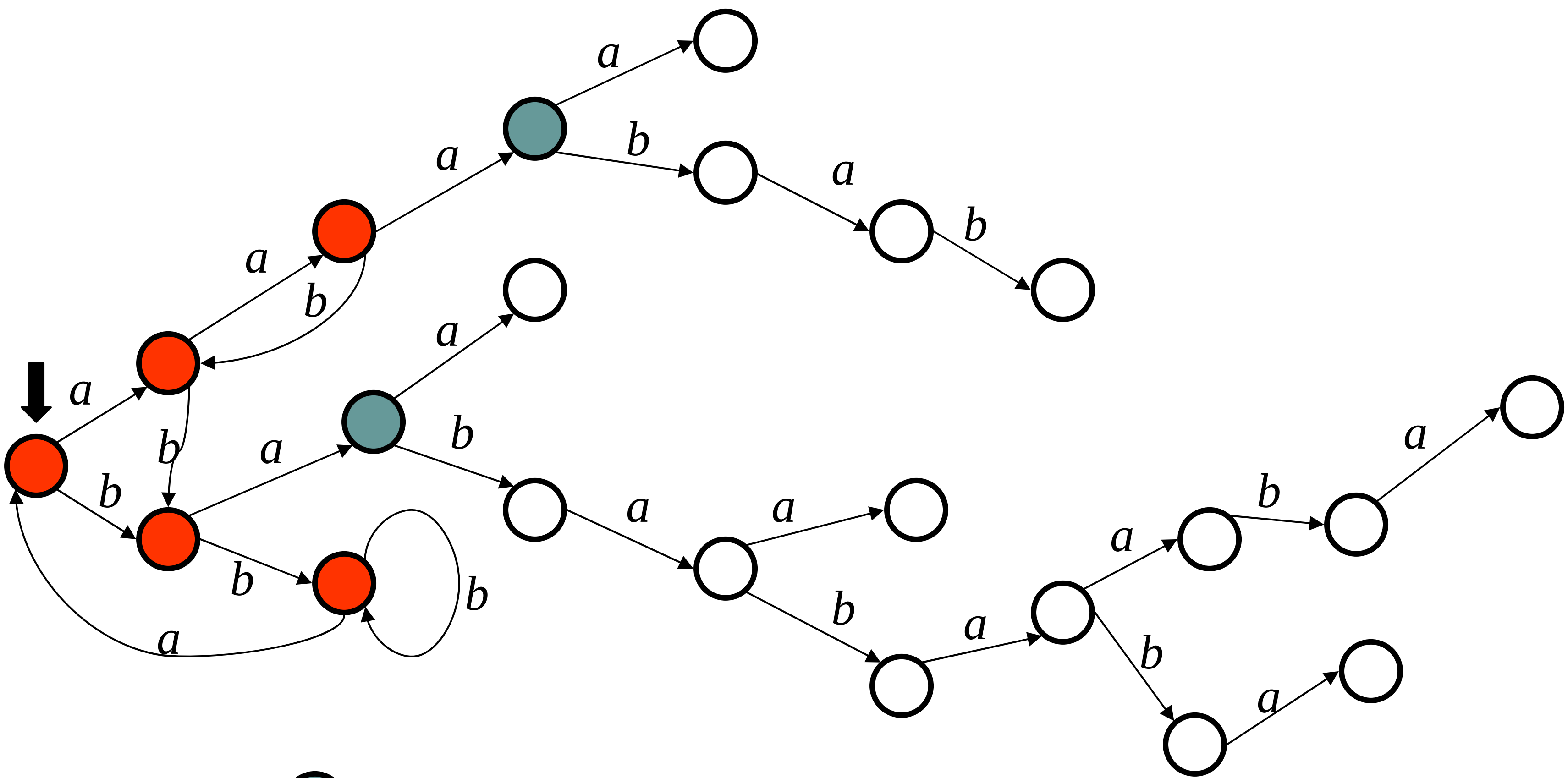
What moves are allowed?

- Merging a with a
- Promoting a to and all its successors that are not to

- Promotion:
- when a can be merged with no

What if there are many merges possible?

- Heuristics
- compute a score
- choose highest score



The blue fringe (Lang 98)

Evidence driven (Lang 98)

for each possible pair (,) **do**

- parse S_+ and S_- on A *resulting from the merge*
- assign a score to each state of A according to the sentences that they accept
- **if** there is a conflict: $-\infty$
- **else** the number of sentences accepted
- **sum** over all states $\Rightarrow$ the score of the merge

select the merge with the highest score

Data driven (cdlh, Oncina & Vidal 96)

For every or state in *A* count

$$v_+(q) = \sum_{w \in S_+} \left|\{u \in \mathrm{Pr}\,ef(w) : \delta(q_0, u) = q\}\right|$$

$$v_-(q) = \sum_{w \in S_-} \left|\{u \in \mathrm{Pr}\,ef(w) : \delta(q_0, u) = q\}\right|$$

Choose the pair (,) such that

$$min(v_+(\ ), v_+(\ )) + min(v_-(\ ), v_-(\ ))$$

is maximal

Careful

- Count first…

  …then try to merge

- Keep track of all tries
- if some is not mergeable, promote it!

Main differences

- data driven is cheaper
- evidence driven won Abbadingo competition

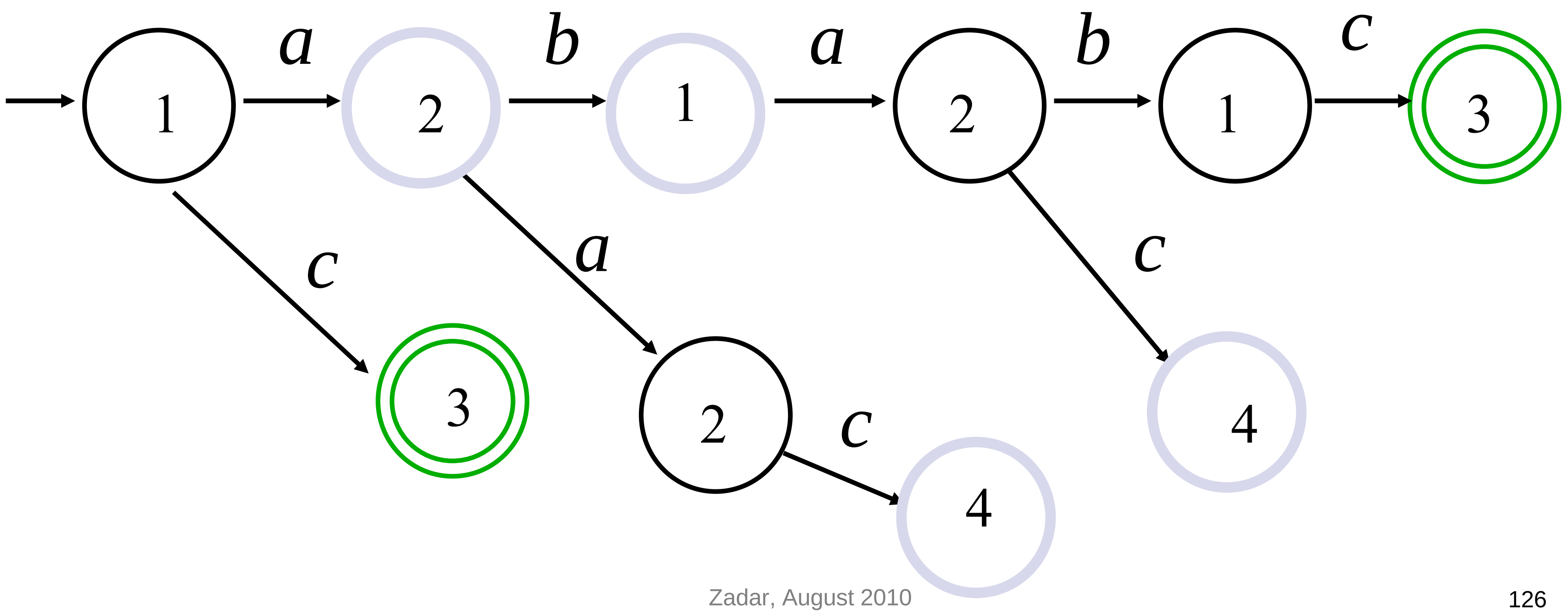
- In the stochastic case, it seems that data driven is a good option…

6.4 Constraint Satisfaction

PTA

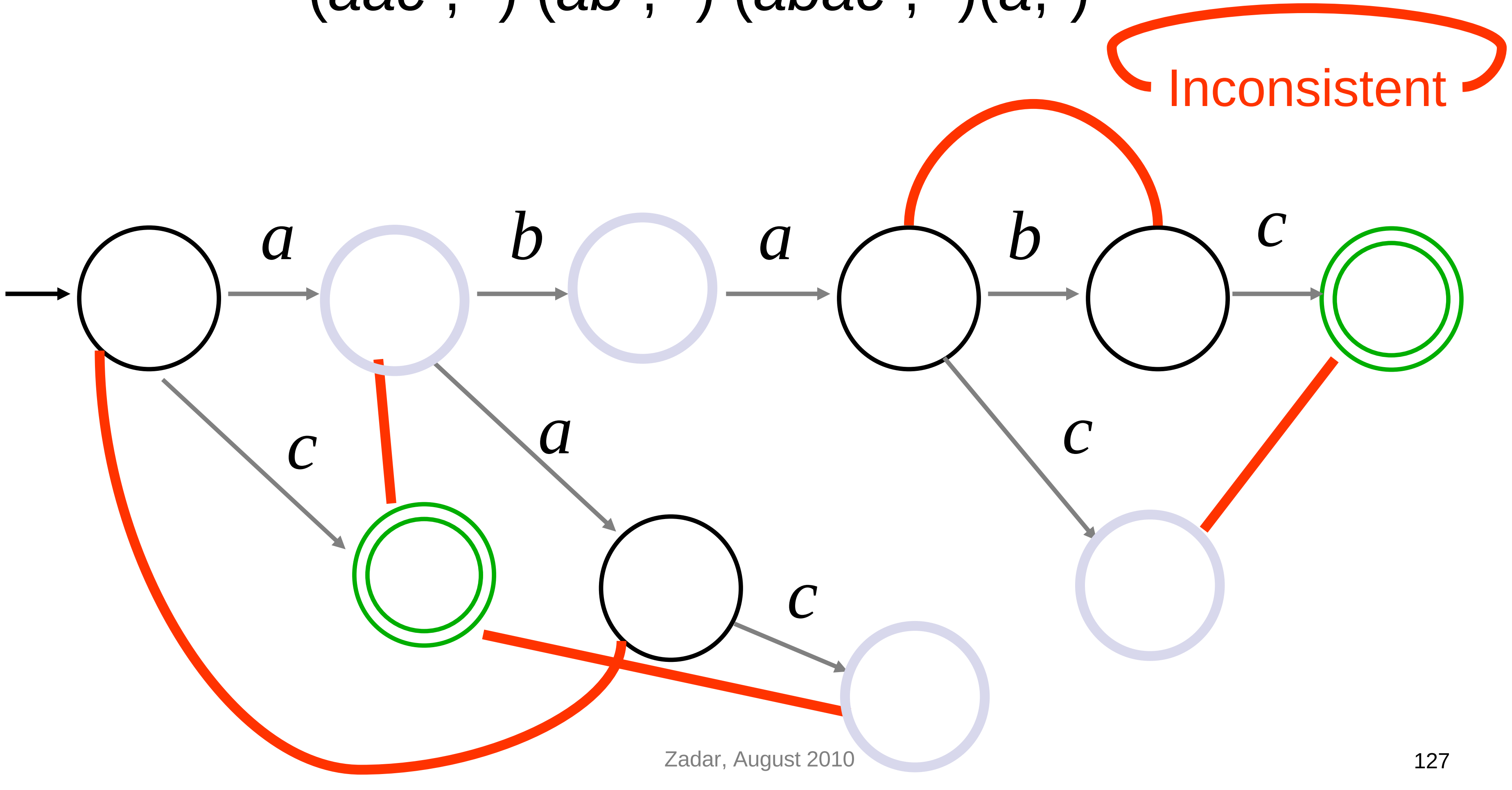
(*ababc* , +) (*c* , +)

(*aac* , -) (*ab* , -) (*abac* , -)(*a*,-)



($ababc$, +) (c , +)

(aac , -) (ab , -) ($abac$, -)(a,-)



Consider (*Q*, *incompatible*)

- All you have to do is find a maximum clique…
- Another NP-hard problem, but for which good heuristics exist.
- Careful : the maximum clique only gives you a lower bound…

Alternatively

- You have $|Q|$ variables $S_1..S_{|Q|}$, and n values $1..n$.
- You have constraints

$$S_i \neq S_j$$

or $$S_i = S_j \Rightarrow S_k = S_l$$

Solve.

Biermann 72, Oliveira & Silva 98, Coste & Nicolas 98

7 Open questions and conclusions

Other versions

- A Matlab version of RPNI

http://www.sec.in.tum.de/~hasan/matlab/gi_toolbox/1.0-Beta/

- A JAVA version
- http://pagesperso.lina.univ-nantes.fr/~cdlh/Downloads/RPNIP.tar.gz

- A parallel version exists, and an OCAML, C, C++…

Some open questions

- Do better than EDSM (still some unsloved Abbadingo task out there…)
- Write a $O(\| f(n) \|)$ algorithm which identifies DFA in the limit (Jose Oncina and cdlh have a log factor still in the way)
- Identify and study the collusion issues
- Deal with noise.